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## Overview

- History and Analysis of Data Hiding Methods
- Exploring NTFS
- FragFS Implementation
- Demonstration
- Detection
- Future Considerations
- A&O

## Why is Data Hiding important?

- Do rootkits alleviate the need for data hiding?
  - NO!!
  - Rootkits hide data on a live system
  - Hiding from forensic/offline analysis is much harder
- What if the rootkit is entirely memory-based?
  - Offline disk analysis will not find it
  - No reboot persistence
- Some techniques rely on covert, persistent storage
  - "The patient hacker"

## History of Data Hiding

- Information Hiding is Old News
  - Writing with invisible ink
- Hiding data on computers is often just a modern day application of existing principles
- Three major categories of data hiding:
  - Out-of-Band
  - In-Band
  - Application Layer



# History of Data Hiding Out-of-Band

### Definition:

- The portion of a medium that is outside the normal specifications for that medium
- The "Media Management" Layer
- Examples:
  - Slack space beyond the end of a partition
  - Slack space at the end of files
    - Example: slacker.exe
  - Sectors marked as bad
  - Host Protected Area



# History of Data Hiding In-Band

### Definition:

- The portion of a medium that is inside the normal specifications for that medium
- The "File System" Layer
- Hidden data must not break the format of the specification

### • Examples:

- Alternative File Streams
- File-System Journal Logs
- Reserved but unallocated sectors



# History of Data Hiding Application Layer

### • Definition:

- Hiding in a higher-level format specification
- Often a subset of In-Band Data Hiding viewed at a different level of granularity

### • Examples:

- Steganography (hiding data within data)
- Hidden text within documents
  - Example: extra white space, tabs, new-line characters
- Virus hiding within EXE's code (.text) section
  - Hydran uses redundancies in i386 code to hide data

## Analysis of Hiding Methods

- Well known to Forensic Tools
  - Forensic tools will specifically look for known hiding methods
    - Alternative File Streams
    - Slack space at the end of files
  - A strings search over a raw disk will find textual results wherever they are located
- Experienced Analysts will detect anomalies not directly identified by Forensic Tools

## Out-of-Band Analysis

- "Coloring Outside the Lines"
- Strengths
  - Being outside the boundaries usually results in being overlooked
  - There is often a large amount of space available
  - Hard to discover without special tools
  - Resilient
- Weaknesses
  - Hard to access without special tools
  - Hard to hide from plain-sight analysis of the outof-band area

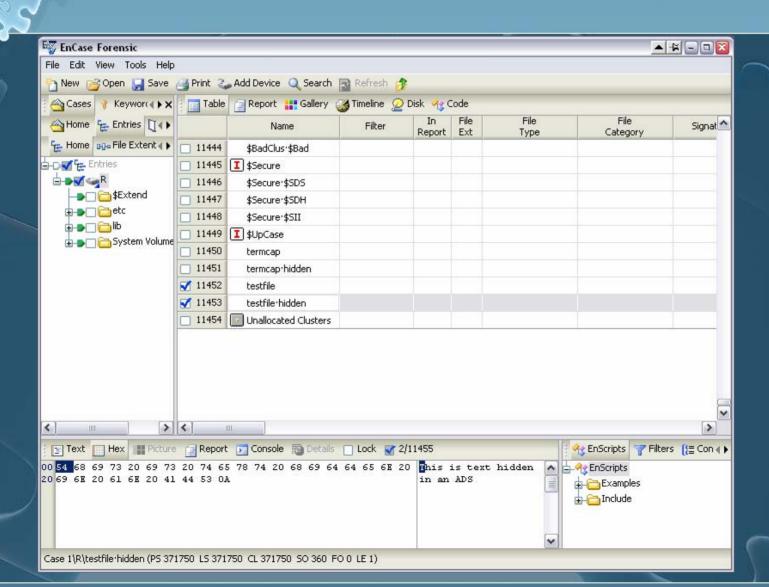
## In-Band Analysis

- "Coloring in the Nooks and Crannies"
- Strengths
  - Usually easy to access with existing tools
  - Follows the specifications
    - Less devious?
- Weaknesses
  - Storage space is often small
  - Relies on security through obscurity easy to detect once method is known
  - Specifications may change

## Application Layer Analysis

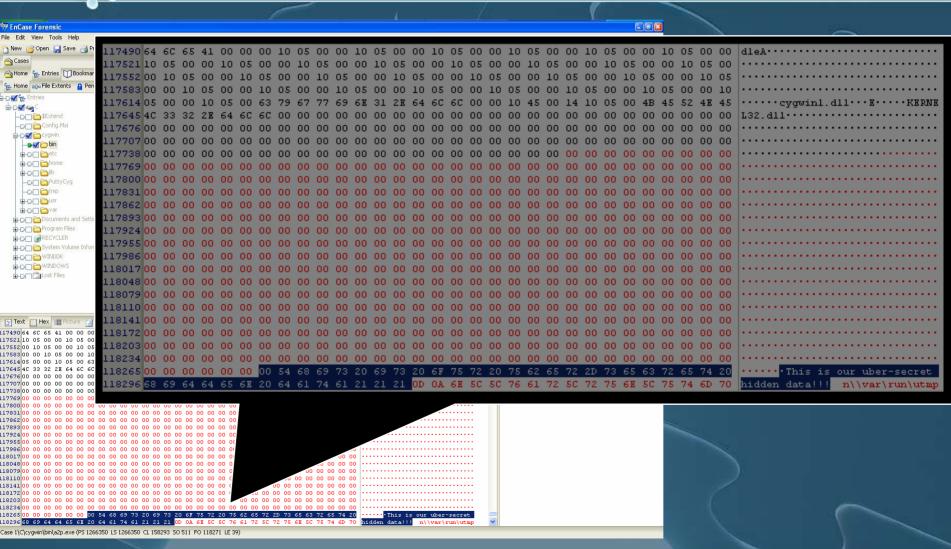
- "Splatter-Painting the Canvas"
- Strengths
  - Hiding in plain sight
  - Often hard to detect
- Weaknesses
  - Storage quantity varies with the size of underlying data, but must be relatively small to remain hidden
  - Difficult to access without special tools
  - Complex algorithms to hide/retrieve data
  - Not resilient

## EnCase – Alternate File Streams





## EnCase - Slacker.exe



## Finding New Places to Hide

- Determine constraints
  - How much space is needed?
  - What type of access is required?
  - How sensitive is hidden data?
- Decide which hiding category best fits the constraints
- Look for previously-unknown hiding methods in that category
  - Analyze an existing specification
    - May require reverse-engineering
  - Study existing data hiding techniques
  - Find unused reserved or slack space

## An NTFS Overview

- Standard file system on Windows NT, Windows 2000, Windows XP, and upcoming Windows Vista
- Master File Table (MFT)
  - Every file or directory is an entry in the table
  - Stores all file system metadata in one place
  - Can grow, but not shrink
  - Not well documented or understood

### MFT Entries

- Each entry is of fixed size
  - Defined in the boot sector
- Each file and directory usually requires one entry but can span multiple entries if needed
- Information about an entry is stored as attributes
  - Each entry has multiple attributes
    - Most files have a few common attributes
    - Attributes can be stored in any order
- Has per sector fix up bytes to detect defects
  - Last two bytes of each sector stored in header and fixed up on every read and write

## MFT Attributes

- Attributes have different types
- Some attribute types can be repeated
  - Duplicate \$DATA attributes commonly called Alternate File Streams
  - Directories entries stored as individual attributes.
- Each attribute can be named, compressed, encrypted, etc
- Each attribute is either resident or nonresident
  - Resident attributes stored within MFT entry
  - Non-resident attributes stored as data run (extent)

## MFT Attribute Examples

- All entries have
  - \$STANDARD\_INFORMATION
    - Stores timestamps, owner ID, security ID, etc
  - \$FILE\_NAME
    - Name by which an entry is known, size, and create/rename timestamp
- All files have \$DATA attribute
- Directories use several attributes-
  - Each entry in a directory is stored as a \$FILE\_NAME attribute
  - DOS 8.3 name stored in a second \$FILE\_NAME attribute
  - Directories have additional indexing attributes to improve filename lookup performance
- End of attributes in an entry is marked by 0xFFFFFFFF
- Most attribute types are kept for backward compatibility



## MFT Entry

#### MFT Entry Header

Magic **US Offset** US Size LSN Sequence Number Hard Links Attribute Offset

Flags

Real Size Allocated Size

Base Entry Number

**Next Attribute ID** 

Reserved

**Entry Number US** Data

#### MFT Entry Header

MFT Attribute

MFT Attribute

MFT Attribute

US

US

Attribute End Mark

Slack Space

#### Resident Attribute

Type Length Resident Flag Name Length Name Offset Flags ID Attribute Length Attribute Offset Reserved Attribute Name

Attribute Data

#### Non-Resident **Attribute**

Type Length Resident Flag Name Length Name Offset Flags ID Data Run Offset Attribute Name Data Run

## Usable Space in MFT entries

- Reserved space within entries
  - Many small unused areas
    - 2 bytes reserved in every entry header
    - 4 byte reserved in resident attributes
    - Up to 14 bytes are reserved in non-resident attributes
    - All attributes are 8 bytes aligned
  - Each file typically has 32 usable bytes
  - Each directory typically has 64 usable bytes
- Slack space after entry attributes
  - Files and directories typically have less than 450 bytes of attributes
  - Default NTFS file systems allocate 1024 bytes per MFT entry
  - Almost 600 bytes per entry!



- Common concerns
  - Entries may be deleted
  - Entries zeroed on allocation
- Reserved Space
  - Might change in future versions of NTFS
  - Normally these bytes are zeroed
- After-attribute slack space
  - Attributes might expand or be added
  - Commonly zero but not always
    - Attributes shrink due to going from resident to nonresident, but can't go back to being resident
    - All directories start as resident and go to non-resident, but can't go back to being resident
    - Attributes can be removed

# Avoiding Pitfalls

- How do we find "safe" entries?
- Many files are rarely modified or deleted
  - Operating system files (drivers, .inf, font, and help files)
  - Most installed application files are only read
  - If it has never been modified it most likely never will be
  - Files that have been around for a long time are rarely deleted
- Non-resident attributes can never become resident
- Directories are rarely deleted
  - Non-resident directories in particular
- Summary Choose entries that are
  - Non-resident
  - Have never been modified

## Putting It All Together

- How much space is available?
  - Base Windows XP Professional install has over 12,000 MFT entries
  - Typical systems have over 100,000 MFT entries
  - Not all entries are safe to use, but testing has shown ~60% of MFT entries are "safe" to use
  - 100,000 entries x 60% x 600 bytes/entry
     36,000,000 bytes!



# Additional Issues Chunking

- Small scattered chunks are not very useful
- The mapping problem
  - Need an interface that can map large blocks or streams across many chunks
  - No matter what space is being used it should look like one contiguous block to higher-level applications
- Mapping should be dynamic
  - Users will delete old files and directories and add new ones
  - Might lose data or need to use additional entries



# Additional Issues Encryption

- Data can be found by searching the raw device
- Detected data can still be protected
- How good is good enough?
  - XOR
  - Blowfish
  - LRW-AES (Narrow-block Encryption)
- Encryption plus Integrity
  - Self-Authenticating encryption is best
- Key management is hard
  - Static forensic analysis can be made difficult
  - Dynamic forensic analysis can always find the keys



# Additional Issues Change Tracking/Redundancy

- What happens when Windows updates an entry you are using?
  - NTFS only changes what it needs to change
  - Might lose some but not all of your data
- Keep extra copies
  - How much redundancy is enough?
- Do your changes get noticed by NTFS?
- Watch for NTFS changing an entry



# Additional Issues Usability

- How is the data presented to the user?
- How is the data presented to the OS?
- Use standard interfaces
  - Prevent the need to rewrite applications
- Reading and writing data files is easy
- Files execution is hard
  - Windows will only execute files from a file system that it understands



## FragFS On-Disk Implementation

- Format
  - Scan MFT Table for suitable entries
    - Non-resident files that have not been modified within the last year
  - Calculate how much space is available in each entry
  - Divide space into 32 byte chunks
    - 4 bytes for Logical Chunk Number
    - 28 bytes for data
- No book marking or index of chunks on disk
  - Check the last 32 bytes of every entry to see if it is a valid chunk
  - If a chunk is valid then check 32 bytes before it the same entry to see if it is also a valid chunk



# FragFS On-Disk Implementation

- Advantages
  - Unlimited redundancy
  - Modification detection
  - Localization of data corruption
  - Easy to relocate or replicate individual chunks of data
- Disadvantages
  - Must scan entire MFT to make updates



## FragFS In-Memory Implementation

- Stackable block device interface
  - Easy to update and add new features
  - On disk format can easily change
- User-space Application Library
  - Can be linked to and used by any application
  - Built-in mini file system
- Kernel Device Driver
  - Creates a virtual disk
  - Uses the FAT file system
  - Can execute files directly from it!

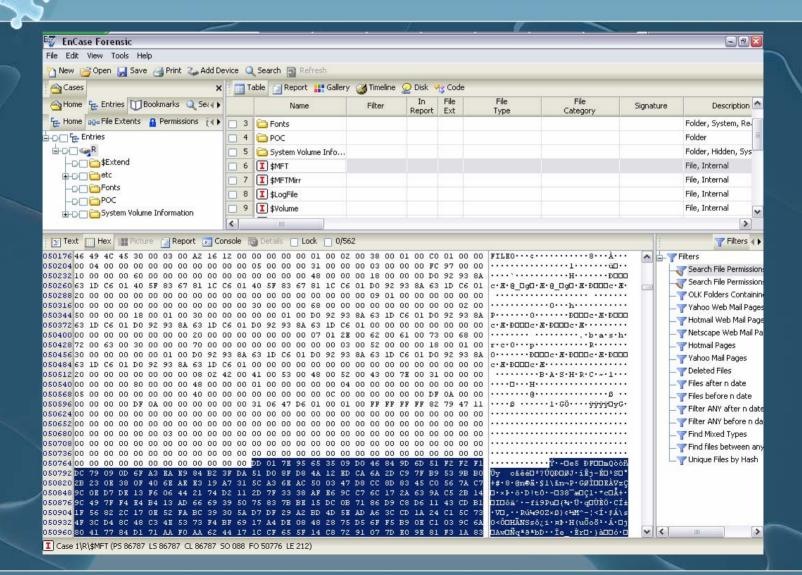


Demonstration

## Detecting NTFS Anomalies

- Current forensic tools treat the MFT as a black box
  - There is a need for forensic tools to better understand file system structures
  - Forensic Analysts do not often have the time to comb through hex dumps
- We have developed a detection tool for data hidden in MFT entry slack space
  - Any data beyond the End-of-Attribute marker is considered suspicious

## Encase - FragFS





### **Future Considerations**

- "Hiding through Obscurity" only buys you time
- Many other unexplored data storage areas
- Hiding access tools is still a problem
  - Bootstrap out of the hidden space?
- Should file system standards be open?
  - Forensic tools could better detect hidden data
  - File systems will be easier to exploit



## **A.8O**



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### Contributors

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