Amendments Proposals to PKCS#11

for support of

WTLS and TLS PRF

This document extends

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PKCS#11 v2.11: Cryptographic	RSA Laboratories November 2001
Token Interface Standard	http://www.rsasecurity.com/rsalabs/PKCS/pkcs-

11/index.html

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1 Introduction

This document contains proposals for amendments to the PKCS#11 Cryptographic Token Interface Standard. The purpose of these amendments is to provide support through a standardized PKCS#11 interface for better secondary PIN handling and to provide support for WTLS. We also propose an amendment for TLS support. New data types and mechanisms are described.

We suggest a standardized way to support WTLS a TLS derived transport security layer that is used in WAP environments.

1.1 Terminology

Definition/Abbreviation	Explanation
IV	Initialization vector
PKCS	Public-Key Cryptography Standards
PRF	Pseudo random function
RSA	The RSA public key crypto system
SW	Software
TBD	To be defined
TLS	Transport Layer Security
WIM	Wireless Identification Module
WTLS	Wireless Transport Layer Security

1.2 References

No	Title	Document No
1	PKCS#11 v2.11: Cryptographic Token Interface	RSA Laboratories November 2001
	Standard	http://www.rsasecurity.com/rsalabs/
		PKCS/pkcs-11/index.html
2	Wireless Transport Layer Security	Wireless Application Protocol
	Version 06-Apr-2001	WAP-261-WTLS-20010406-a
		http://www.wapforum.org/
3	The TLS Protocol Version 1.0	RFC 2246
		The Internet Engineering Task
		Force, January 1999
		http://www.ietf.org/

1.3 Yet to do

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2 New general data types

This chapter contains additions to chapter 9 of [1].

2.1 New object types

This chapter contains additions to Chapter 9.4 of [1].

The following additional certificate type is defined.

#define CKC_WTLS TBD

2.2 New data types for mechanisms

This chapter contains additions to Chapter 9.5 of [1].

The following additional mechanism types are defined.

```
#define CKM_WTLS_PRE_MASTER_KEY_GEN TBD
#define CKM_WTLS_MASTER_KEY_DERIVE TBD
#define CKM_WTLS_MASTER_KEY_DERVIE_DH_ECC TBD
#define CKM_WTLS_PRF TBD
#define CKM_WTLS_SERVER_KEY_AND_MAC_DERIVE TBD
#define CKM_WTLS_CLIENT_KEY_AND_MAC_DERIVE TBD
#define CKM_TLS_PRF TBD
```

3 New objects

This chapter contains additions to Chapter 10 of [1].

3.1 New certificate objects

This chapter contains additions to Chapter 10.6 of [1].

The following figure illustrates details of certificate objects and replaces figure 7 of [1]:

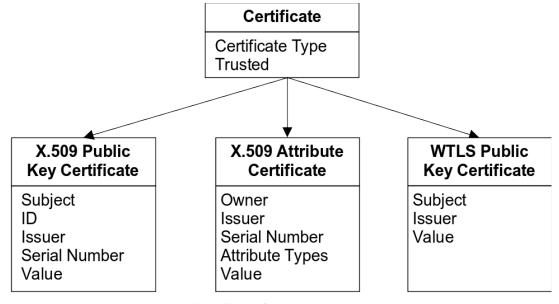


Figure 1, Certificate Object Attribute Hierarchy

3.1.1 WTLS public key certificate objects

Details can be found in [2].

WTLS certificate objects (certificate type **CKC_WTLS**) hold WTLS public key certificates. The following table defines the WTLS certificate object attributes, in addition to the common attributes listed in Table 15 of [1], Table 19 of [1] and Table 21 of [1]:

Table: WTLS Certificate Object Attributes

Attribute	Data type	Meaning
CKA_SUBJECT ¹	Byte array	WTLS-encoding of the certificate
		subject name.
CKA_ISSUER	Byte array	WTLS-encoding of the certificate
		issuer name. (default empty)
CKA_VALUE ¹	Byte array	WTLS-encoding of the certificate.

¹Must be specified when the object is created.

Only the **CKA ISSUER** attribute may be modified after the object is created.

The following is a sample template for creating a certificate object:

4 New mechanisms

This chapter contains additions to Chapter 12 of [1].

4.1 TLS mechanism parameters

Details can be found in [3].

4.1.1 CK_TLS_PRF_PARAMS

CK_TLS_PRF_PARAMS is a structure, which provides the parameters to the **CKM TLS PRF** mechanism. It is defined as follows:

```
typedef struct
{
   CK_BYTE_PTR     pSeed;
   CK_ULONG     ulSeedLen;
   CK_BYTE_PTR     pLabel;
   CK_ULONG     ulLabelLen;
   CK_BYTE_PTR     pOutput
   CK_ULONG_PTR     pulOutputLen;
} CK TLS PRF PARAMS;
```

The fields of the structure have the following meanings:

pSeed pointer to the input seed

ulSeedLen length in bytes of the input seed

pLabel pointer to the identifying label

ulLabelLen length in bytes of the identifying label

pOutput pointer receiving the output of the operation

pulOutputLen pointer to the length in bytes that the output to be

created shall have, has to hold the desired length as input and will receive the calculated length as

output

CK_TLS_PRF_PARAMS_PTR is a pointer to a CK_TLS_PRF_PARAMS.

4.2 TLS mechanisms

Details can be found in [3].

4.2.1 PRF (pseudo random function)

PRF (pseudo random function) in TLS, denoted **CKM_TLS_PRF**, is a mechanism used to produce a secure digest protected by a secret key. It is used to produce a securely generated random output of arbitrary length. The keys it uses are generic secret keys.

It has a parameter, a **CK_TLS_PRF_PARAMS** structure, which allows for the passing of the input seed and its length, the passing of an identifying label and its length and the passing of the length of the output to the token and for receiving the output.

This mechanism produces securely generated random output of the length specified in the parameter.

This mechanism departs from the other key derivation mechanisms in Cryptoki in not using the template sent along with this mechanism during a **C_DeriveKey** function call, which means the template shall be a NULL_PTR, and its returned information. For most key-derivation mechanisms, **C_DeriveKey** returns a single key handle as a result of a successful completion. However, since the **CKM_TLS_PRF** mechanism returns the requested number of output bytes in the **CK_TLS_PRF_PARAMS** structure specified as the mechanism parameter, the parameter *phKey* passed to **C_DeriveKey** is unnecessary, and should be a NULL_PTR.

If a call to **C_DeriveKey** with this mechanism fails, then no output will be generated.

4.3 WTLS mechanism parameters

Details can be found in [2].

4.3.1 CK WTLS RANDOM DATA

CK_WTLS_RANDOM_DATA is a structure, which provides information about the random data of a client and a server in a WTLS context. This structure is used by the CKM WTLS MASTER KEY DERIVE mechanism. It is defined as follows:

```
typedef struct
{
   CK_BYTE_PTR pClientRandom;
   CK ULONG ulClientRandomLen;
```

4.3.2 CK_WTLS_MASTER_KEY_DERIVE_PARAMS

CK_WTLS_MASTER_KEY_DERIVE_PARAMS is a structure, which provides the parameters to the **CKM_WTLS_MASTER_KEY_DERIVE** mechanism. It is defined as follows:

be used (possible types can be found in [2])

RandomInfo clients and servers random data information

pVersion pointer to a **CK_BYTE** which receives the WTLS protocol version information

CK_WTLS_MASTER_KEY_DERIVE_PARAMS_PTR is a pointer to a CK_WTLS_MASTER_KEY_DERIVE_PARAMS.

4.3.3 CK_WTLS_PRF_PARAMS

CK_WTLS_PRF_PARAMS is a structure, which provides the parameters to the **CKM WTLS PRF** mechanism. It is defined as follows:

The fields of the structure have the following meanings:

DigestMechanism the mechanism type of the digest mechanism to

```
be used (possible types can be found in [2])
```

pSeed pointer to the input seed

ulSeedLen length in bytes of the input seed

pLabel pointer to the identifying label

ulLabelLen length in bytes of the identifying label

pOutput pointer receiving the output of the operation

pulOutputLen pointer to the length in bytes that the output to be

created shall have, has to hold the desired length as input and will receive the calculated length as

output

CK_WTLS_PRF_PARAMS_PTR is a pointer to a CK_WTLS_PRF_PARAMS.

4.3.4 CK WTLS KEY MAT OUT

CK_WTLS_KEY_MAT_OUT is a structure that contains the resulting key handles and initialization vectors after performing a C_DeriveKey function with the

CKM_WTLS_SEVER_KEY_AND_MAC_DERIVE or with the

CKM_WTLS_CLIENT_KEY_AND_MAC_DERIVE mechanism. It is defined as follows:

```
typedef struct
{
   CK_OBJECT_HANDLE hMacSecret;
   CK_OBJECT_HANDLE hKey;
   CK_BYTE_PTR pIV;
} CK_WTLS_KEY_MAT_OUT;
```

The fields of the structure have the following meanings:

hMacSecret Key handle for the resulting MAC secret key

hKey Key handle for the resulting secret key

pIV Pointer to a location which receives the initialisation vector (IV) created (if any)

CK_WTLS_KEY_MAT_OUT _PTR is a pointer to a CK_WTLS_KEY_MAT_OUT.

4.3.5 CK_WTLS_KEY_MAT_PARAMS

CK_WTLS_KEY_MAT_PARAMS is a structure that provides the parameters to the CKM_WTLS_SEVER_KEY_AND_MAC_DERIVE and the CKM_WTLS_CLIENT_KEY_AND_MAC_DERIVE mechanisms. It is defined as follows:

CK_WTLS_RANDOM_DATA RandomInfo;
CK_WTLS_KEY_MAT_OUT_PTR pReturnedKeyMaterial;
} CK_WTLS_KEY_MAT_PARAMS;

The fields of the structure have the following meanings:

DigestMechanism the mechanism type of the digest mechanism to

be used (possible types can be found in [2])

ulMacSizeInBits the length (in bits) of the MACing keys agreed

upon during the protocol handshake phase

ulKeySizeInBits the length (in bits) of the secret keys agreed upon

during the handshake phase

ulIVSizeInBits the length (in bits) of the IV agreed upon during

the handshake phase. If no IV is required, the

length should be set to 0.

ulSequenceNumber The current sequence number used for records

sent by the client and server respectively

blsExport a boolean value which indicates whether the keys

have to be derived for an export version of the protocol. If this value is true (i.e. the keys are exportable) then *ulKeySizeInBits* is the length of the key in bits before expansion. The length of the key after expansion is determined by the information found in the template sent along with this mechanism during a C_DeriveKey function call (either the CKA_KEY_TYPE or

the CKA_VALUE_LEN attribute).

RandomInfo client's and server's random data information

pReturnedKeyMaterial points to a CK_WTLS_KEY_MAT_OUT

structure which receives the handles for the keys

generated and the IVs

CK_WTLS_KEY_MAT_PARAMS_PTR is a pointer to a CK_WTLS_KEY_MAT_PARAMS.

4.4 WTLS mechanisms

Details can be found in [2].

4.4.1 Pre master secret key generation for RSA key exchange suite

Pre master secret key generation for the RSA key exchange suite in WTLS denoted CKM_WTLS_PRE_MASTER_KEY_GEN, is a mechanism, which generates a variable length secret key. It is used to produce the pre master secret key for RSA key exchange suite used in WTLS. This mechanism returns a handle to the pre master secret key.

It has one parameter, a CK BYTE, which provides the client's WTLS version.

The mechanism contributes the CKA_CLASS, CKA_KEY_TYPE and CKA_VALUE attributes to the new key (as well as the CKA_VALUE_LEN attribute, if it is not supplied in the template). Other attributes may be specified in the template, or else are assigned default values.

The template sent along with this mechanism during a C_GenerateKey call may indicate that the object class is CKO_SECRET_KEY, the key type is CKK_GENERIC_SECRET, and the CKA_VALUE_LEN attribute indicates the length of the pre master secret key.

For this mechanism, the ulMinKeySize field of the **CK_MECHANISM_INFO** structure indicate 20 bytes.

4.4.2 Master secret key derivation

Master secret derivation in WTLS, denoted **CKM_WTLS_MASTER_KEY_DERIVE**, is a mechanism used to derive a 20 byte generic secret key from variable length secret key. It is used to produce the master secret key used in WTLS from the pre master secret key. This mechanism returns the value of the client version, which is built into the pre master secret key as well as a handle to the derived master secret key.

It has a parameter, a **CK_WTLS_MASTER_KEY_DERIVE_PARAMS** structure, which allows for passing the mechanism type of the digest mechanism to be used as well as the passing of random data to the token as well as the returning of the protocol version number which is part of the pre master secret key.

The mechanism contributes the CKA_CLASS, CKA_KEY_TYPE, and CKA_VALUE attributes to the new key (as well as the CKA_VALUE_LEN attribute, if it is not supplied in the template). Other attributes may be specified in the template, or else are assigned default values.

The template sent along with this mechanism during a **C_DeriveKey** call may indicate that the object class is **CKO_SECRET_KEY**, the key type is **CKK_GENERIC_SECRET**, and the **CKA_VALUE_LEN** attribute has value 20. However, since these facts are all implicit in the mechanism, there is no need to specify any of them.

This mechanism has the following rules about key sensitivity and extractability:

The **CKA_SENSITIVE** and **CKA_EXTRACTABLE** attributes in the template for the new key can both be specified to be either TRUE or FALSE. If omitted, these attributes each take on some default value.

If the base key has its CKA_ALWAYS_SENSITIVE attribute set to FALSE, then the derived key will as well. If the base key has its CKA_ALWAYS_SENSITIVE attribute set to TRUE, then the derived key has its CKA_ALWAYS_SENSITIVE attribute set to the same value as its CKA_SENSITIVE attribute.

Similarly, if the base key has its **CKA_NEVER_EXTRACTABLE** attribute set to FALSE, then the derived key will, too. If the base key has its **CKA_NEVER_EXTRACTABLE** attribute set to TRUE, then the derived key has its **CKA_NEVER_EXTRACTABLE** attribute set to the *opposite* value from its **CKA_EXTRACTABLE** attribute.

For this mechanism, the ulMinKeySize and ulMaxKeySize fields of the **CK MECHANISM INFO** structure both indicate 20 bytes.

Note that the **CK_BYTE** pointed to by the **CK_WTLS_MASTER_KEY_DERIVE_PARAMS** structure's *pVersion* field will be modified by the **C_DeriveKey** call. In particular, when the call returns, this byte will hold the WTLS version associated with the supplied pre master secret key.

Note that this mechanism is only useable for key exchange suites that use a 20-byte pre master secret key with an embedded version number. This includes the RSA key exchange suites, but excludes the Diffie-Hellman and Elliptic Curve Cryptography key exchange suites.

4.4.3 Master secret key derivation for Diffie-Hellman and Elliptic Curve Cryptography

Master secret derivation for Diffie-Hellman and Elliptic Curve Cryptography in WTLS, denoted CKM_WTLS_MASTER_KEY_DERIVE_DH_ECC, is a mechanism used to derive a 20 byte generic secret key from variable length secret key. It is used to produce the master secret key used in WTLS from the pre master secret key. This mechanism returns a handle to the derived master secret key.

It has a parameter, a **CK_WTLS_MASTER_KEY_DERIVE_PARAMS** structure, which allows for the passing of the mechanism type of the digest mechanism to be used as well as random data to the token. The *pVersion* field of the structure must be set to NULL_PTR since the version number is not embedded in the pre master secret key as it is for RSA-like key exchange suites.

The mechanism contributes the CKA_CLASS, CKA_KEY_TYPE, and CKA_VALUE attributes to the new key (as well as the CKA_VALUE_LEN attribute, if it is not supplied in the template). Other attributes may be specified in the template, or else are assigned default values.

The template sent along with this mechanism during a **C_DeriveKey** call may indicate that the object class is **CKO_SECRET_KEY**, the key type is **CKK_GENERIC_SECRET**, and the **CKA_VALUE_LEN** attribute has value 20. However, since these facts are all implicit in the mechanism, there is no need to specify any of them.

This mechanism has the following rules about key sensitivity and extractability:

The **CKA_SENSITIVE** and **CKA_EXTRACTABLE** attributes in the template for the new key can both be specified to be either TRUE or FALSE. If omitted, these attributes each take on some default value

If the base key has its **CKA_ALWAYS_SENSITIVE** attribute set to FALSE, then the derived key will as well. If the base key has its **CKA_ALWAYS_SENSITIVE** attribute set to TRUE, then the derived key has its **CKA_ALWAYS_SENSITIVE** attribute set to the same value as its **CKA_SENSITIVE** attribute.

Similarly, if the base key has its **CKA_NEVER_EXTRACTABLE** attribute set to FALSE, then the derived key will, too. If the base key has its **CKA_NEVER_EXTRACTABLE** attribute set to TRUE, then the derived key has its **CKA_NEVER_EXTRACTABLE** attribute set to the *opposite* value from its **CKA_EXTRACTABLE** attribute.

For this mechanism, the ulMinKeySize and ulMaxKeySize fields of the **CK_MECHANISM_INFO** structure both indicate 20 bytes.

Note that this mechanism is only useable for key exchange suites that do not use a fixed length 20-byte pre master secret key with an embedded version number. This includes the Diffie-Hellman and Elliptic Curve Cryptography key exchange suites, but excludes the RSA key exchange suites.

4.4.4 PRF (pseudo random function)

PRF (pseudo random function) in WTLS, denoted **CKM_WTLS_PRF**, is a mechanism used to produce a secure digest protected by a secret key. It is used to produce a securely generated random output of arbitrary length. The keys it uses are generic secret keys.

It has a parameter, a **CK_WTLS_PRF_PARAMS** structure, which allows for passing the mechanism type of the digest mechanism to be used, the passing of the input seed and its length, the passing of an identifying label and its length and the passing of the length of the output to the token and for receiving the output.

This mechanism produces securely generated random output of the length specified in the parameter.

This mechanism departs from the other key derivation mechanisms in Cryptoki in not using the template sent along with this mechanism during a **C_DeriveKey** function call, which means the template shall be a NULL_PTR, and its returned information. For most key-derivation mechanisms, **C_DeriveKey** returns a single key handle as a result of a successful completion. However, since the **CKM_WTLS_PRF** mechanism returns the requested number of output bytes in the **CK_WTLS_PRF_PARAMS** structure specified as the mechanism parameter, the parameter *phKey* passed to **C_DeriveKey** is unnecessary, and should be a NULL_PTR.

If a call to C **DeriveKey** with this mechanism fails, then no output will be generated.

4.4.5 Server Key and MAC derivation

Server key, MAC and IV derivation in WTLS, denoted

CKM_WTLS_SERVER_KEY_AND_MAC_DERIVE, is a mechanism used to derive the appropriate cryptographic keying material used by a cipher suite from the master secret key and random data. This mechanism returns the key handles for the keys generated in the process, as well as the IV created.

It has a parameter, a **CK_WTLS_KEY_MAT_PARAMS** structure, which allows for the passing of the mechanism type of the digest mechanism to be used as well as random data as well as the characteristic of the cryptographic material for the given cipher suite and a pointer to a structure which receives the handles and IV which were generated. This structure is defined in Section 4.3.4

This mechanism contributes to the creation of two distinct keys on the token and returns one IV (if an IV is requested by the caller) back to the caller. The keys are all given an object class of **CKO_SECRET_KEY**.

The MACing key (server write MAC secret) is always given a type of **CKK_GENERIC_SECRET**. It is flagged as valid for signing, verification and derivation operations.

The other key (server write key) is typed according to information found in the template sent along with this mechanism during a **C_DeriveKey** function call. By default, it is flagged as valid for encryption, decryption, and derivation operations.

An IV (server write IV) will be generated and returned if the *ulIVSizeInBits* field of the **CK_WTLS_KEY_MAT_PARAMS** field has a nonzero value. If it is generated, its length in bits will agree with the value in the *ulIVSizeInBits* field

Both keys inherit the values of the CKA_SENSITIVE, CKA_ALWAYS_SENSITIVE, CKA_EXTRACTABLE, and CKA_NEVER_EXTRACTABLE attributes from the base key. The template provided to C_DeriveKey may not specify values for any of these attributes that differ from those held by the base key.

Note that the CK_WTLS_KEY_MAT_OUT structure pointed to by the CK_WTLS_KEY_MAT_PARAMS structure's *pReturnedKeyMaterial* field will be modified by the C_DeriveKey call. In particular, the two key handle fields in the CK_WTLS_KEY_MAT_OUT structure will be modified to hold handles to the newly-created keys; in addition, the buffer pointed to by the CK_WTLS_KEY_MAT_OUT structure's *pIV* field will have the IV returned in them (if an IV is requested by the caller). Therefore, this field must point to a buffer with sufficient space to hold any IV that will be returned.

This mechanism departs from the other key derivation mechanisms in Cryptoki in its returned information. For most key-derivation mechanisms, **C_DeriveKey** returns a single key handle as a result of a successful completion. However, since the

CKM_WTLS_SERVER_KEY_AND_MAC_DERIVE mechanism returns all of its key handles in the CK_WTLS_KEY_MAT_OUT structure pointed to by the CK_WTLS_KEY_MAT_PARAMS structure specified as the mechanism parameter, the parameter *phKey* passed to C DeriveKey is unnecessary, and should be a NULL PTR.

If a call to **C_DeriveKey** with this mechanism fails, then *none* of the two keys will be created on the token.

4.4.6 Client key and MAC derivation

Client key, MAC and IV derivation in WTLS, denoted

CKM_WTLS_CLIENT_KEY_AND_MAC_DERIVE, is a mechanism used to derive the appropriate cryptographic keying material used by a cipher suite from the master secret key and random data. This mechanism returns the key handles for the keys generated in the process, as well as the IV created.

For this mechanism all applies as described in the Chapter except for that the names server write MAC secret, server write key and server write IV have to be replaced by client write MAC secret, client write key and client write IV.

REMARK: When comparing the existing TLS mechanisms in Cryptoki with these extensions to support WTLS one could argue that there would be no need to have distinct handling of the client and server side of the handshake. However since in WTLS the server and client have different sequence numbers for the server and the client. There could be instances where WTLS is used to protect asynchronous protocols and where sequence numbers on the client and server side therefore would not be necessarily aligned, and hence this motivates the introduced split..