Introduction to Public-Key Technology

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Outline

Concepts

Secret key, public key, message digest, ...

Algorithms

DES, RSA, DSS, ...

RSA Details

Applications

Digital signatures and envelopes

Standards

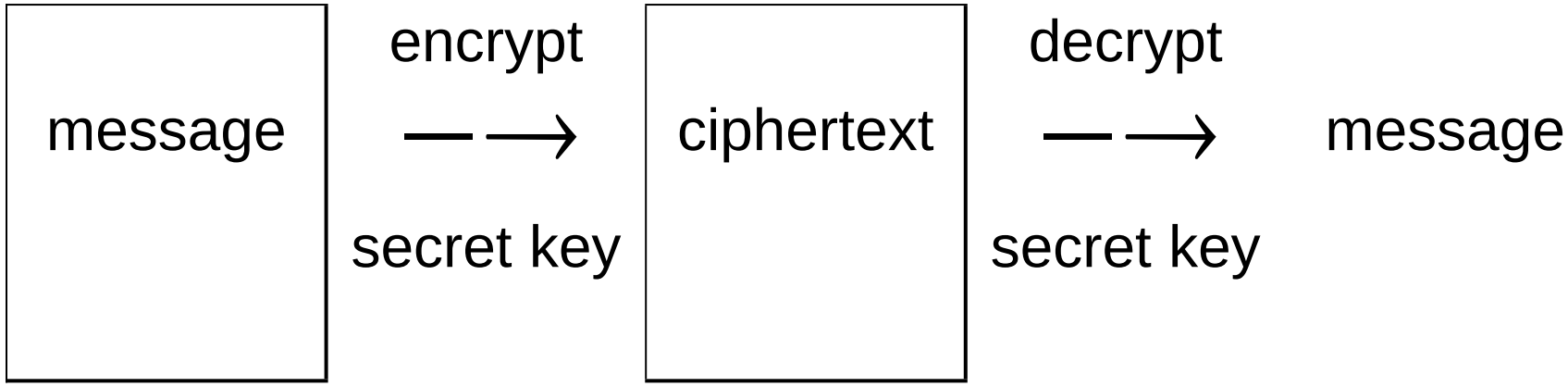
Conclusions

Concepts

Secret-key cryptosystem

Encryption, decryption with same key.

For privacy.



Users agree secretly on key.

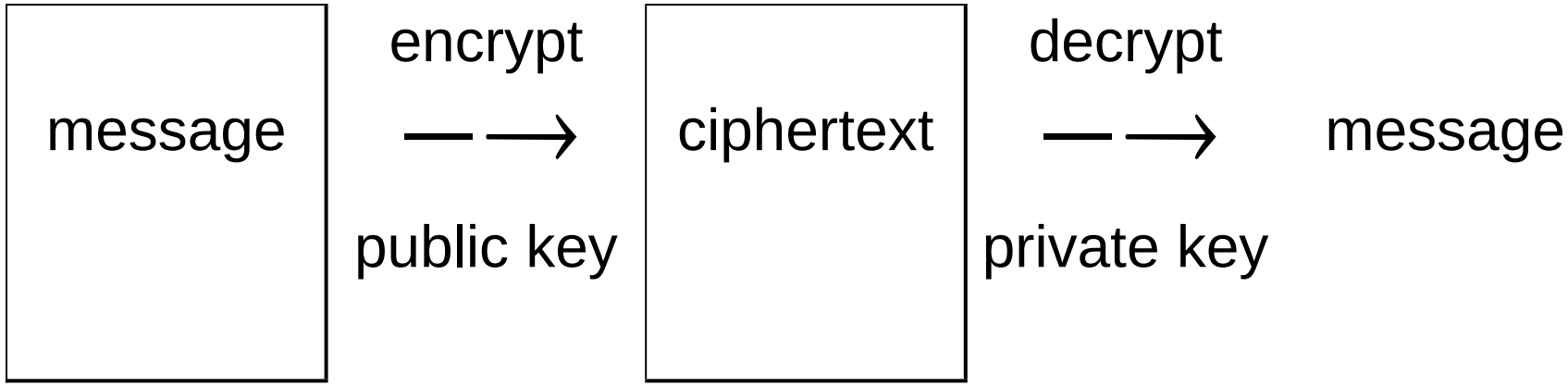
Examples: DES, RC2, RC4.

Concepts (cont'd)

Public-key cryptosystem

Encryption, decryption with different keys.

For privacy.



Users keep one key private, publish other.

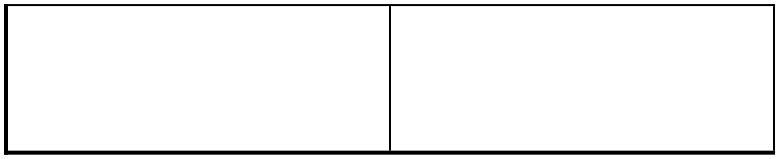
Examples: RSA, ElGamal.

Often hybrid with secret key.

Concepts (cont'd)

Secret key vs. public key

	secret key	**public key**
copies/ secret	two	one
secrets/user	many	one
scalability	fair	good
speed	good	fair

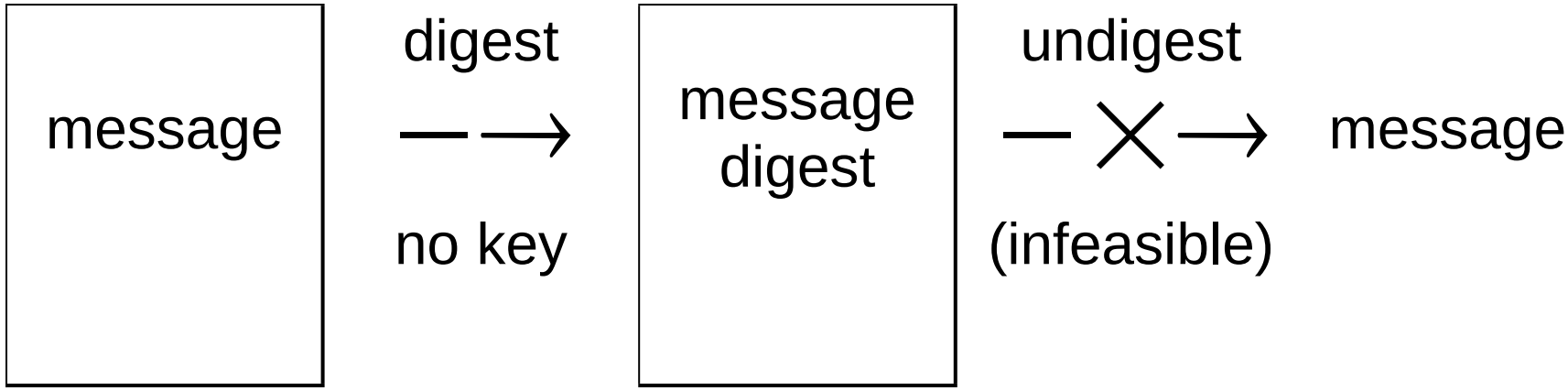


Hybrid cryptography combines benefits.

Concepts (cont'd)

Message-digest algorithm

For "fingerprinting"—one-way hash.



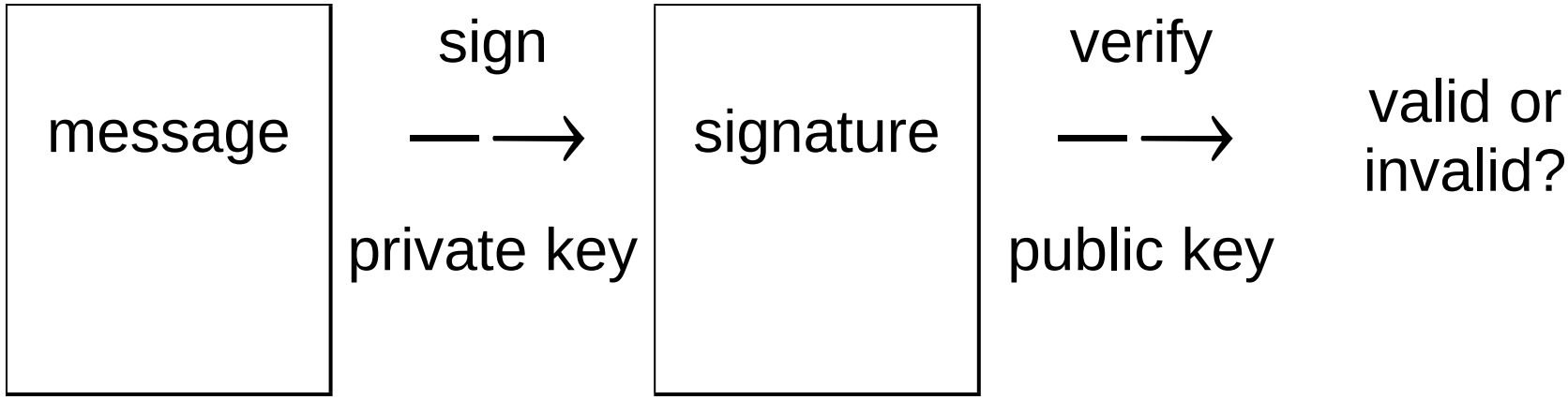
Typically 128 or 160 bits.

Examples: MD5, SHS.

Concepts (cont'd)

Digital signature scheme

- Signature with private key, verification with public.
- For authentication—of message and signer.



Examples: RSA, DSS.

Usually hybrid with message digest.

Algorithms

Data Encryption Standard (DES)

- NBS, 1976
- Secret-key cryptosystem
- 56-bit keys

RSA

- Rivest-Shamir-Adleman, MIT, 1977
- Public-key cryptosystem *and* digital signature scheme
- Based on difficulty of factoring large integers

Algorithms (cont'd)

RC2, RC4

- Rivest, RSADSI, 1980s
- Block & high-speed stream secret-key cryptosystems
- Variable-length keys

MD5

- Rivest, MIT/RSADSI, 1991
- High-speed message-digest algorithm
- 128-bit digest

Algorithms (cont'd)

Digital Signature Standard (DSS)

- NIST, 1991
- Digital signature scheme
- Based on difficulty of discrete logarithms

Secure Hash Standard (SHS)

- NIST, 1992
- High-speed, DSS-compatible message-digest algorithm
- 160-bit digest

RSA Details

Keys

n: public modulus

e: public exponent (typically 3 or $2^{16}+1$)

d: private exponent

p,*q*: private factors of modulus

$$n = p \times q$$

$$d \times e \bmod (p\text{-}1)(q\text{-}1) = 1$$

Public key is (*n*,*e*).

Private key is (*n*,*d*).

RSA Details (cont'd)

Encryption with public key

m: message

Ciphertext = c where

$$c = m^e \bmod n.$$

m may be a key.

Decryption with private key

$$m = c^d \bmod n.$$

RSA Details (cont'd)

Encryption with private key

Signature = s where

$$s = m^d \bmod n.$$

m may be a message digest.

Decryption with public key

$$m = s^e \bmod n.$$

RSA Details (cont'd)

Performance

RSA operations involve *modular multiplication*, which takes time proportional to $(\log n)^2$.

Public-key: 2 to 17 multiplications

Private-key: 1.5 log *n* multiplications

Given *p*, *q*, four times faster.

Good public-key speed, fair private-key speed
—but good in combination with secret key, message digest.

RSA Details (cont'd)

Performance (cont'd)

Examples with 512-bit keys, $e = 2^{16}+1$, given p, q:

Processor	Public key (sec.)	Private key (sec.)
16 MHz 68020	.32	3.3
12 MHz 80286	.25	2.7
25 MHz 68040	.065	.65
20 MHz 80386	.065	.55

30 MHz DSP16A	.035	.17
20MHz DSP56000	.0081	.044

For 1024 bits, public key $\times$ 4, private key $\times$ 8.

RSA Details (cont'd)

Security

Goal: Given n, find p and q.

Typical approaches take time

$$L(n) = \exp((1+\varepsilon)).$$

$L(n)$ is *subexponential* in $\log_2 n$:

For any constant c, $L(n)$ grows slower than n^c (exponential), faster than $(\log_2 n)^c$ (polynomial).

Thus, hardware speedups help multiplication more than factoring.

RSA Details (cont'd)

Security (cont'd)

Based on $L(p)$ as instruction count (Rivest, 1991):

$\log_2 n$	$L(n)$	MIPS years
512	6.7×10^{19}	2.1×10^{6}
576	1.7×10^{21}	5.5×10^{7}
	. . .	
960	3.7×10^{28}	1.2×10^{15}

1024	4.4×10^{29}	1.4×10^{16}

MIPS year = one million instructions/second for one year = 3.1×10^{13} intructions.

$2^{56} \approx 7.2 \times 10^{16}$ (not directly comparable).

Applications

Hybrid cryptography

Digest + public key = digital signature.

Secret key + public key = digital envelope.

Performance, scalability, no shared secrets.

Two tools

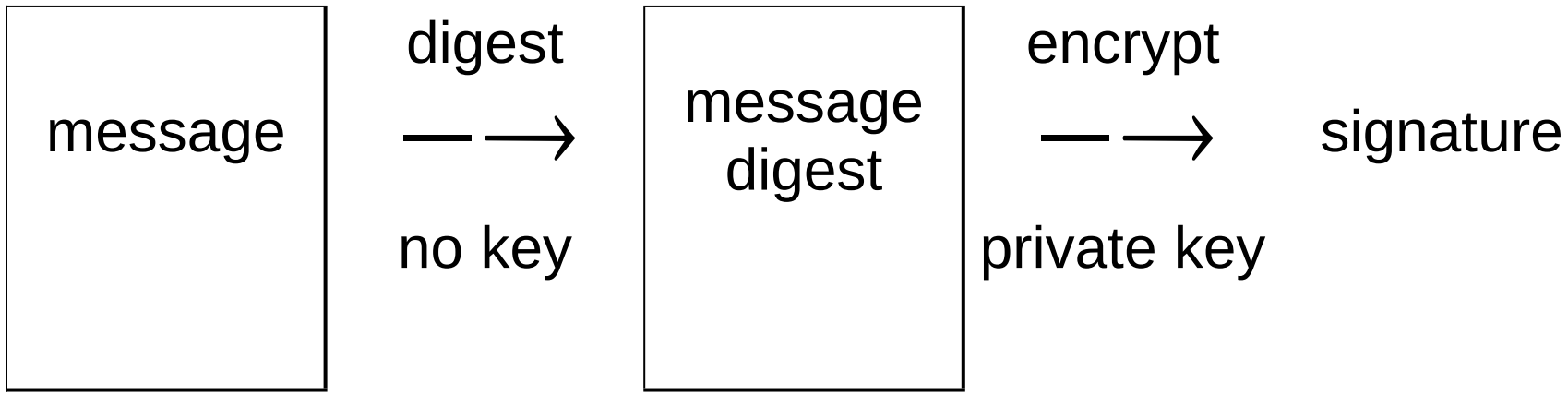
Digital signatures: *sign* and *verify*.

Digital envelopes: *seal* and *open*.

Applications (cont'd)

Signing a message

Alice digests message, encrypts signature with her private key.

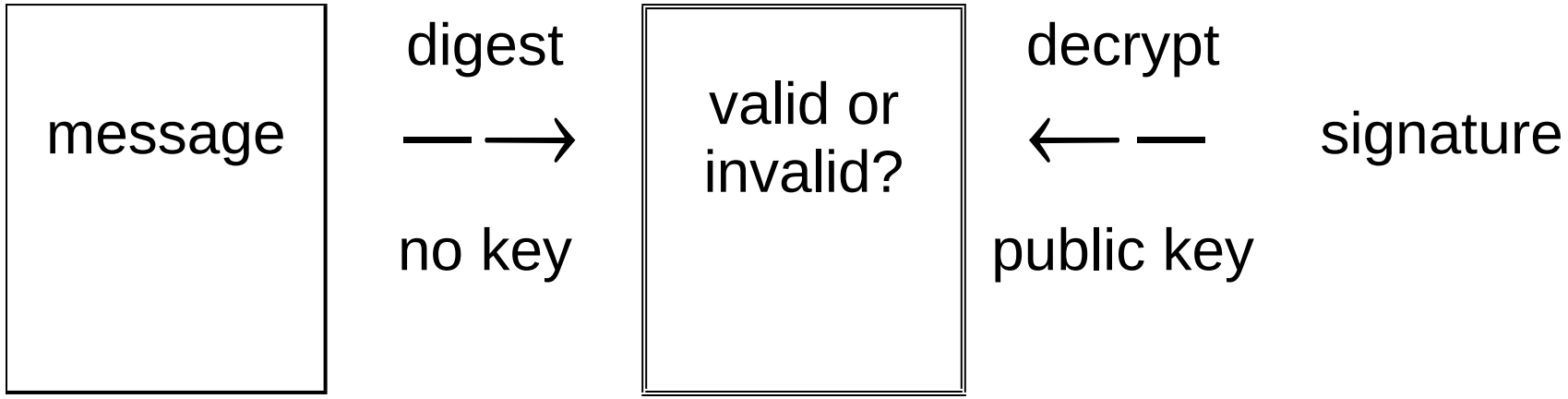


She sends the message and signature to Bob.

Applications (cont'd)

Verifying a signature

Bob digests message, decrypts signature with Alice's public key, compares results.

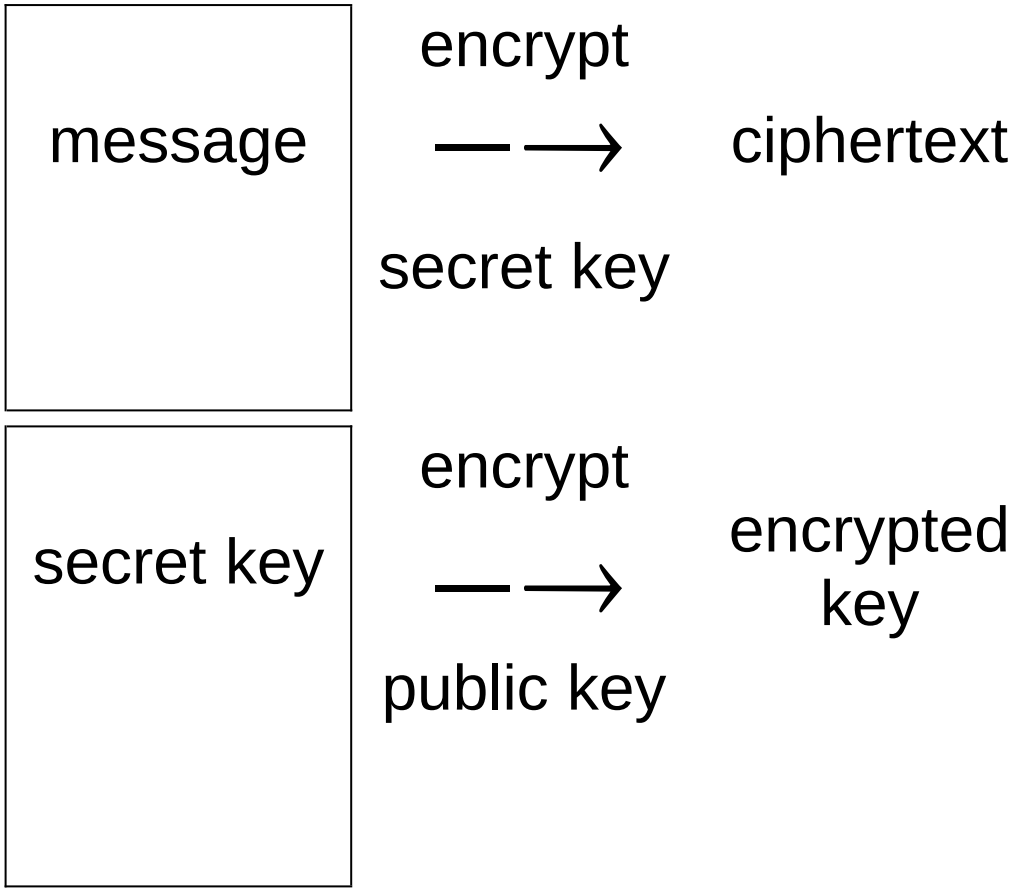


The signature is valid if and only if the results are the same.

Applications (cont'd)

Sealing a message

Alice encrypts message with a secret key,
encrypts secret key with Bob's public key.

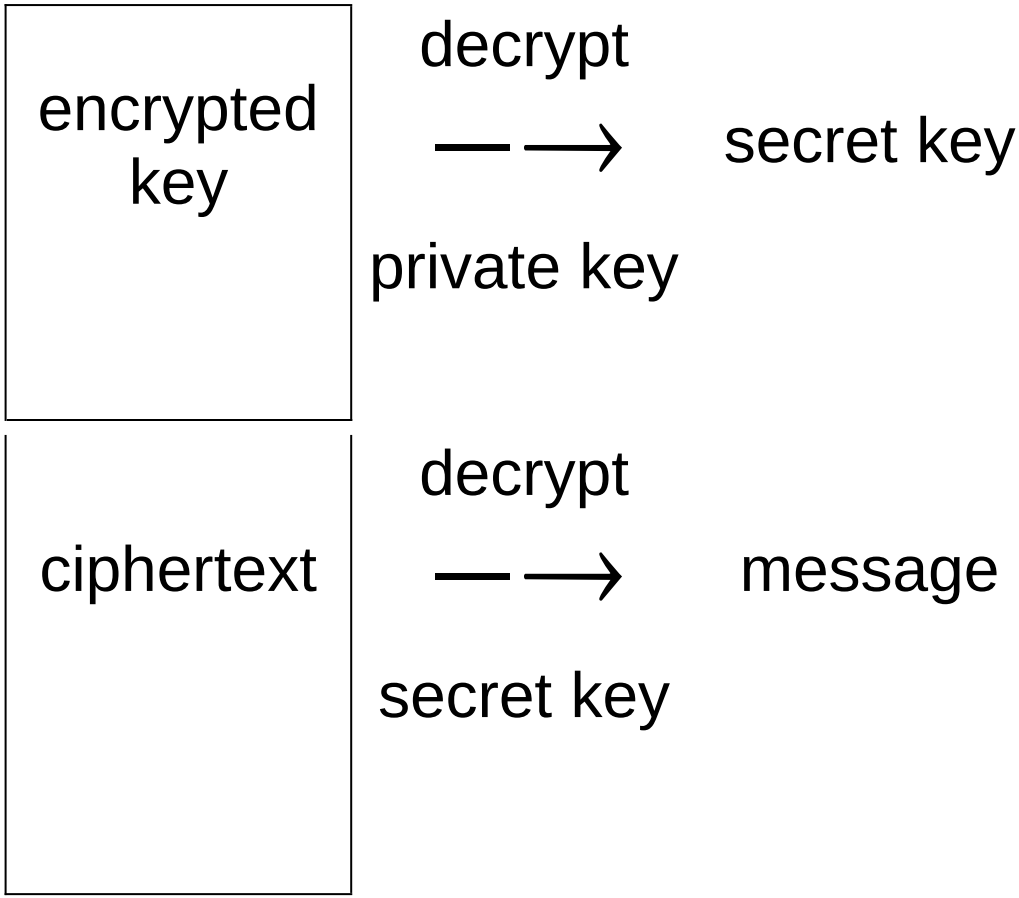


She sends the ciphertext and encrypted key to Bob.

Applications (cont'd)

Opening an envelope

Bob decrypts the encrypted key with his private key, decrypts the ciphertext with the secret key.


encrypted key
decrypt
secret key
private key
ciphertext
decrypt
message
secret key

Standards

ANSI

X9.30,.31: DSS, RSA

NIST

DES, DSS, SHS

Key management, certification forthcoming

CCITT

X.400, X.500

ISO/IEC

IS 9796: RSA-oriented signature scheme

Standards (cont'd)

Internet

Privacy-Enhanced Mail: RSA, DES, MD5

SNMP: MD5

RSADSI *et al*

PKCS (Public-Key Cryptography Standards)

Also: **French Banking, Standards Australia, ...**

Conclusions

Basic concepts, various algorithms, many standards

Secret key, public key, message digest

RSA, DES, MD5, DSS, ...

ANSI, NIST, ISO, PKCS, ...

Powerful applications

Digital signatures, digital envelopes

Hybrid cryptography combines benefits