

# Overview and History of Operating Systems

**These are the notes for lecture 1.**

**Please review the  
“Syllabus” notes  
before these.**

# **Overview / Historical Developments**

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## **An Operating System...**

**Sits between hardware and users**

**Provides “environment” to execute programs**

**Like a government**

**No useful work**

**Regulates workers**

**Manages, allocates resources**

**CPU (execution time)**

**Memory Space**

**Disk / File storage**

**I/O Devices**

**Control**

**Prevent incorrect use of hardware**

**Security / Protection**

# Goals

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Make computer *easy* to use

Make computer more *efficient*

Help user *solve problems / do work*

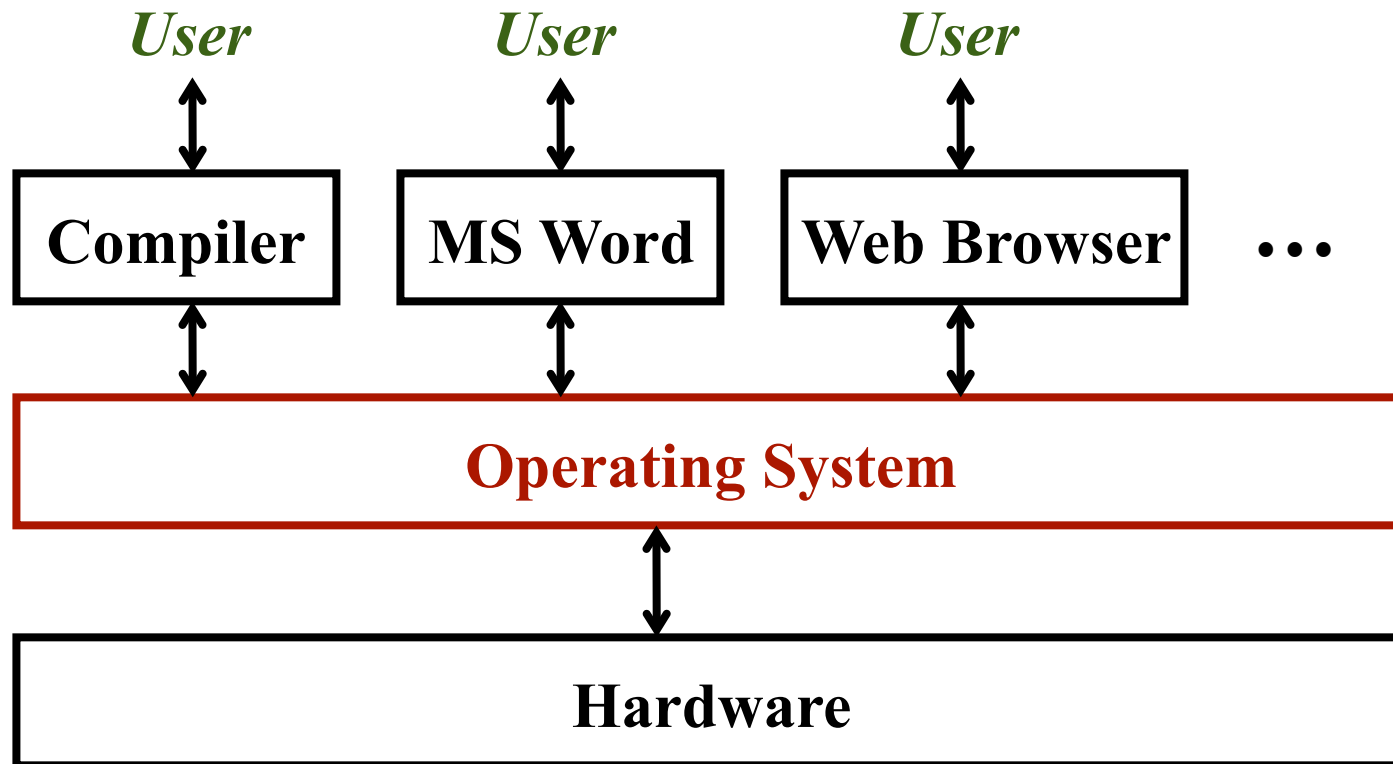
Ease of use

Efficiency

} *Often in conflict*

# Overview

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# Early Computers

## Input Devices:

- Card Reader

## Output Devices:

- Printer
- Card Puncher



Punch card

**(No disk, no secondary storage)**

A “*job*”

User prepares input cards (Program, Data)

User gets time on the machine

Loads the program

Executes it

Study the output and come back tomorrow

# Early Computers

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## The O.S....

(Simply a “*control program*”)

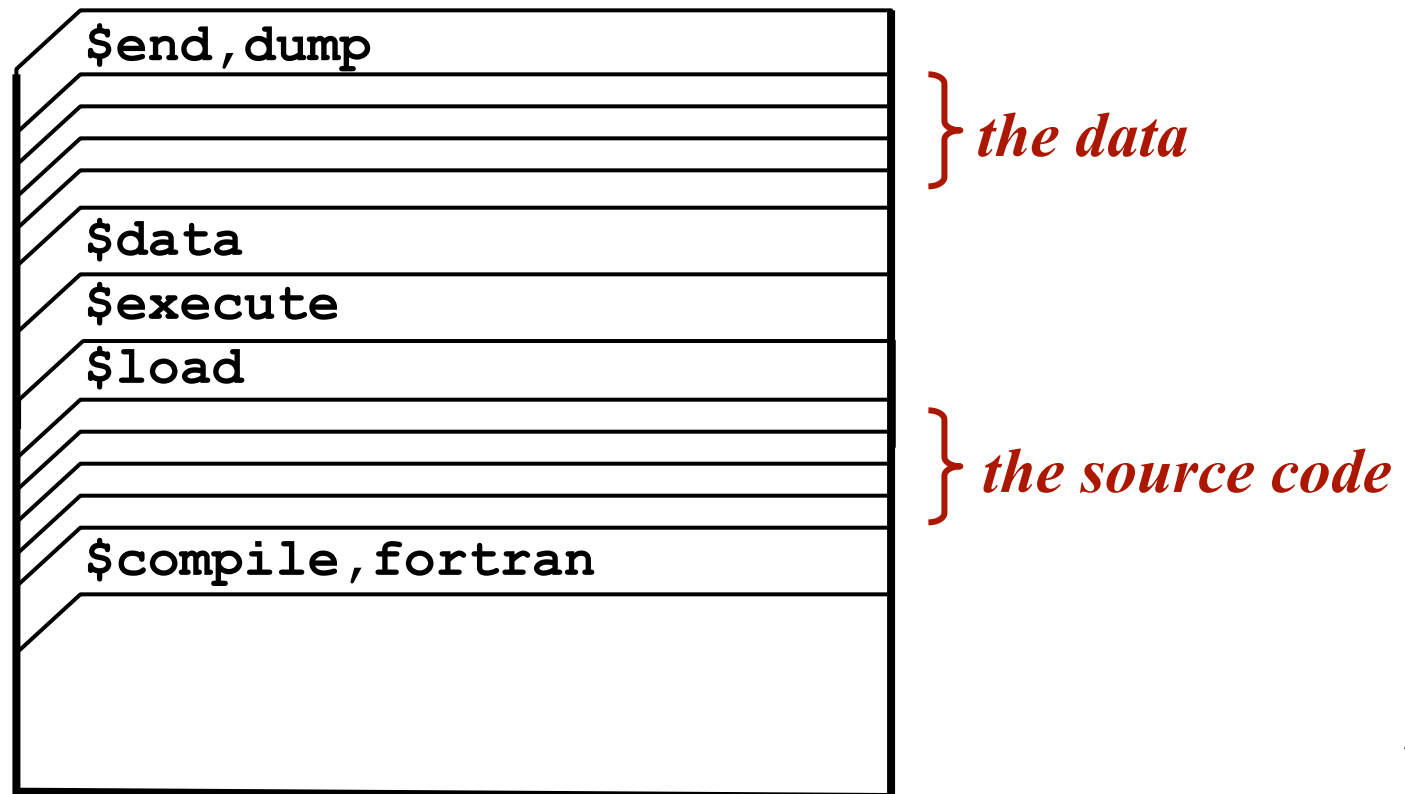
- Read cards
- Load memory
- Transfer control to user code
- Print out contents of memory (“*core dump*”)
- Loop to next “*job*”

# Batch Operating System

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Read jobs from cards

Some jobs are “control cards”



# **New Technology: Magnetic Tape**

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**Idea:** Read cards onto tape

To output a card...

Write to tape & punch it later

Same for printing



# New Technology: Magnetic Tape

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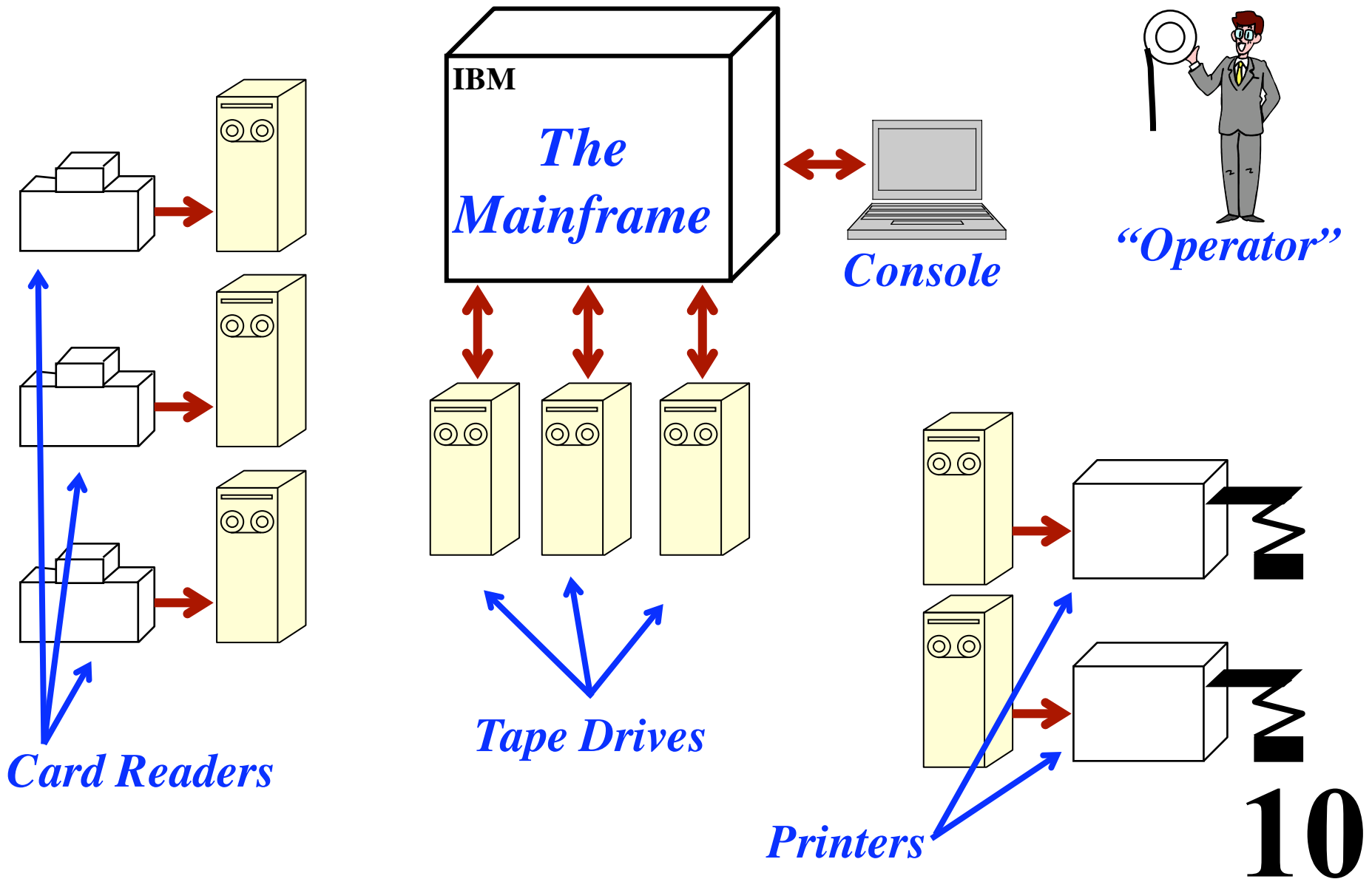
Same for printing

Concept: ***SPOOLING***

- Use tape as a “buffer”.
- *Allows I/O from one job to overlap the computation from another job!*

(When disks were invented, they were first used like this.)

# A Mainframe System (circa 1960)



# **New Technology: The Disk**

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**The first disks were used for spooling.  
(The “file” was invented later.)**

**Concept: The Job Pool**

**Several jobs are waiting to be executed**

- **One job in memory**
- **Future jobs sitting on disk**

# **New Technology: The Disk**

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**Concept: The Job Pool**

**Several jobs are waiting to be executed**

- **One job in memory**
- **Future jobs sitting on disk**

**The O.S. can make its first decision!**

***Which job to run next?***

**First-come, first served**

**vs.**

**Job Scheduling**

# Multi-programming

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*Idea: Keep several jobs in memory at once!*

When one job waits on I/O...  
another jobs can use the CPU!

**Increases CPU utilization**

**Don't keep entire job pool in memory  
(just select a few)**

**Job X starts I/O...**

**OS selects another job to run.**

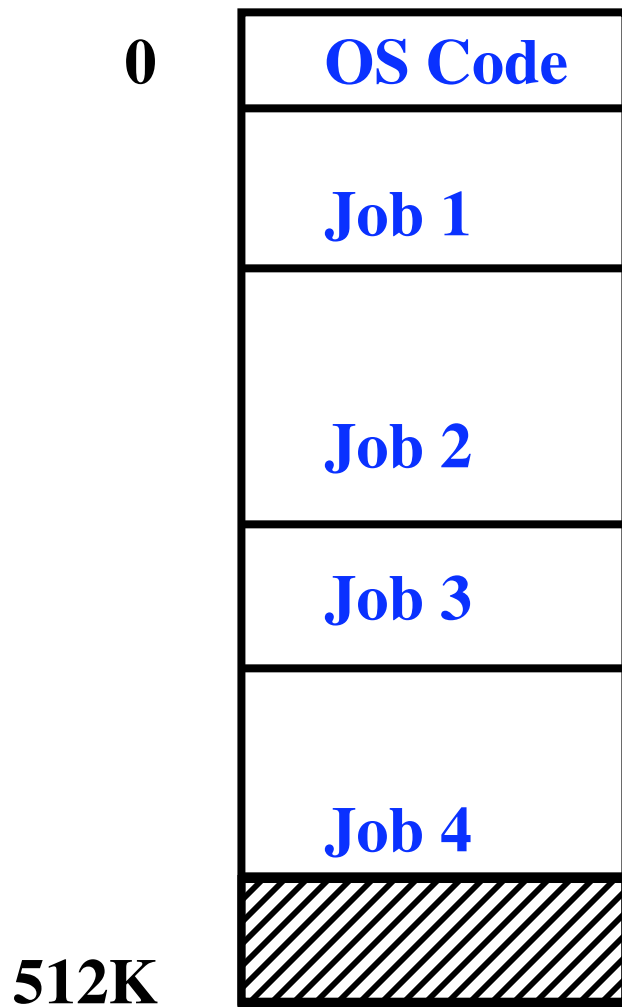
**CPU does not sit idle.**

**When job X finishes...**

**OS selects another job from job pool  
and loads it into memory.**

# Main Memory

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# Main Memory

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## Simplest Approach:

- The entire job is loaded into a contiguous range of memory
- No protection between programs
- A job runs until it requests I/O  
No “*time-slicing*”  
If it loops... Oh, well...
- When complete, it transfers back to the OS.

## Memory Management

A big topic (a chapter in textbook)

## Job Scheduling

When job X blocks (due to I/O)...

... which job will be run next?

“Process Management” (chapter 2)

# **New Concept: Terminals**

**Combines with multiprogramming!**

**Users want to interact with the programs  
*while they are executing!***

**Initial motivations for interaction:**

- **Deal with contingencies during job execution.**
- **Debugging programs.**

***The good 'ole days...***

- **Long turn-around times.**
- **Program were smaller.**
- **Written in assembly / FORTRAN.**
- **Programmers were very careful.**
- **Bugs were intolerable.**



# **Batch Operating Systems**

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**Each job runs non-stop  
(until it decides to perform I/O)**

**Problems when used with interactive terminals:**

- **User response times are...**
  - Long**
  - Unpredictable**
- **Bugs in one program...**
  - Crash the entire system**
  - Really annoyed people using terminals**
  - (With batch jobs, just restart the programs)**

# **New Concept: Time-Sharing**

**“Multi-tasking”**

## **Goal:**

**Make it seem like every user has a dedicated computer!**

## **The Idea:**

**Switch the CPU rapidly between jobs so that every running job seems to be making regular progress.**

## **New hardware support required:**

**Periodic “timer interrupts”**

**Ways to protect one program from the other programs.**

**Goal: isolate bugs / loops to just that program!**

# Historical Context

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**Each user has a Cathode Ray Tube (CRT) terminal.**

**Each user types a command & it is executed.**

**“Response time” -- should be < 1 second.**

**Many users “online”, sharing the CPU.**

**Demonstrated in 1960, common in 1970s.**

**Files kept on disks...**

**but lots of implementation details still visible.**

**(blocking, file formats, sectors, etc...)**

**UNIX** (1969, spreading in mid-1970's)

# **Time-Sharing**

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**Several processes in memory.**

**Memory management & protection are required.**

**If size of user job is really large...**

**Swap other users' jobs out to disk (temporarily)**

***“Virtual Memory”***

**Don't load the entire job into memory**

## **Disk Management:**

**Directory Structures**

**Ease-of-use**

**Protection from other users**

## **Communication between users:**

**Still an active area**

**Synchronization, coordination**

# **Personal Computers**

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**Began appearing in the 1970s.**

**Shift of emphasis:**

- **Hardware utilization is less important**
- **Maximize: User convenience**
  - Fast response time**
  - Ease of programming**
- **Security vs. Communication**

**Early opinions: Protection is unimportant**

**Each person's computer is separate & isolated.**

# Trends

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Hardware costs will continue to fall.

*OS's will appear embedded in more devices!*

*More kinds of OS's will be needed!*

Computer users will become more sophisticated.

*Better OS's will be demanded!*

Features on research / high-end systems will become common on small, inexpensive systems.

*Parallel processing / multi-processing*

*Real-time control*

*Distributed systems*

Programs will become more complex.

*The OS will need to promote ease of programming / use!*

More malicious programs will be created.

*More security will be needed!*

# **Kinds of Operating Systems**

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**Personal Computers**

**Large, super-computers**

**Parallel multi-processors**

**Embedded computers**

**Real-time computers**

**Distributed systems**

**Highly reliable systems**

**Super-low cost**

**Research platforms**

## **Applications**

- **Chess playing**
- **Dishwasher / microwave**
- **Flight control / space shuttle**
- **Military command**
- **Automobile control**
- **Lab / factory automation**
- **Assembly robot control**
- **Corporation management**
- **Web server**
- **Web search engine**
- **Nuclear reactor control**
- **Toys**
- **Artificial Intelligence**

*... etc ...*

# Parallel Systems

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## Single processor vs. multi-processor systems

### Goals:

- Increase *“Throughput”*
  - Get more work done, per hour
- Utilize small, inexpensive processors
  - ... To get more horsepower (giga-flops)
  - Example: Graphics co-processor
- Save money by sharing expensive peripheral devices.
- Reliability
  - Graceful degradation
  - Fault-tolerance



# Fault-Tolerance

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## Example: Tandem System

Two identical processors

- Primary
- Backup

Each has its own memory

Operating in lock-step

Failure detected?

Backup becomes the primary

*“hot backup”*

# **Types of Multiprocessing Systems**

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## ***Tightly-Coupled Systems***

- **Share the system bus**
- **Share peripheral I/O devices**
- **Share memory (sometimes)**
- **Share a common clock (sometimes)**

**Example: Graphics Co-processor**

## ***Lossely-Coupled Systems***

# **Types of Tightly-Coupled Systems**

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## **“Symmetric Multiprocessing”**

- **Each processor runs identical copy of OS**
- **OS code resides in shared memory**
- **Shared data structures (for concurrency control)**

## **“Asymmetric Multiprocessing”**

- **Master-slave relationship**
- **One processor assigns tasks to others**
- **Increased specialization**
  - > **decreased reliability**
- **Trend: cheap processors**
  - Offload tasks to slaves or back-ends**
    - Graphics processor**
    - Disk processor**
    - Processor in keyboard / mouse**
- **Communication/interfaces becomes paramount.**

# **Distributed Systems**

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**Processors do not share**

**Memory, Clock, System bus, Devices**

**“Loosely-coupled systems”**

**Communication**

**Slow (internet)  $\longleftrightarrow$  Fast (specialized bus)**

**Motivations**

**Communication & sharing**

**Email, shared data, group work**

**Reliability**

**No single failure should crash the entire system**

**Data should remain accessible**

**Computation speed-up**

**Load-sharing**

**Resource sharing**

**Access to specialized or unique hardware**

# Real-Time Operating Systems

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Used to control a physical process

Sensors collect data (“input”?)

Actuators control the physical process (“output”?)

## Examples:

- Avionics (space shuttle control system)
- Medical systems (heart monitor)
- Industrial system (oil refinery)
- Military (anti-missile laser control)
- Consumer products (automobile controller)

The key...

Rigid, well-defined, fixed time constraints

Time-sharing systems *should* respond quickly!

Real-time systems *must* respond quickly!

# **Real-Time Operating Systems**

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## **Soft Real-Time Systems**

**Some processes are given higher “*priorities.*”**

**Example: video & music playback**

**Common in many OS’s.**

**Adequate for many applications,**

**... but too risky for some!**

**Even though response is often fast enough,**

**it is never guaranteed!**

# Real-Time Operating Systems

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## Hard Real-Time Systems

Guarantees that task will complete by their deadlines.

All potential delays are *bounded*.

Want to avoid:

- **Disks** (highly variable latencies)
- **Virtual Memory** (complex & unpredictable)

OS tends to be low-level, minimal, close to hardware.

A specialized sub-field of OS.