#### **Chapter 2 (Fourth Part)**

# Scheduling Mechanisms and Policies

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#### **Process state model**



**CPU Utilization** – *How busy is the CPU*?

**Throughput** – *How many jobs finished (per unit time)*?

**Turnaround Time** – *How long from job submission to job termination*?

**Response Time** – How long (on average) does it take to get a "response" from a "stimulus"?

**Missed deadlines** – Were any deadlines missed?

#### **Scheduler Options**

#### **Priorities**

#### May use priorities to determine who runs next amount of memory, order of arrival, etc..

#### **Dynamic vs. Static algorithms**

Dynamically alter the priority of the tasks while they are in the system (possibly with feedback)

Static algorithms typically assign a fixed priority when the job is initially started.

#### **Preemptive vs. Nonpreemptive**

Preemptive systems allow the task to be interrupted at any time so that the O.S. can take over again.

#### **Scheduling Policies**

First-Come, First Served (FIFO) Shortest Job First (non-premeptive) Shortest Job First (with preemption) Round-Robin Scheduling Priority Scheduling Real-Time Scheduling

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	Arrival	Processing	
Process	Time	Time	
1	0	3	
2	2	6	
3	4	4	
4	6	5	
5	8	2	



Process	Arrival Time	Processing Time		Turnaround
			Delay	Time
1	0	3		
2	2	6		
3	4	4		
4	6	5		
5	8	2		

Total time taken,









Total time taken, From submission to







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**Goal:** Enable interactivity

Limit the amount of CPU that a process can have at one time.

*Time quantum* 

Amount of time the OS gives a process before intervention The "time slice" Typically: 1 to 100ms

	Arrival	Processing
Process	Time	Time
1	0	3
2	2	6
3	4	4
4	6	5
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	Arrival	Processing	
Process	Time	Time	
1	0	3	
2	2	6	
3	4	4	
4	6	5	
5	8	2	



		Arrival	Processing		Turnaround
Proc	cess	Time	Time	Delay	Time
	1	0	3	1	4
	2	2	6	10	16
	3	4	4	9	13
	4	6	5	9	14
	5	8	2	5	7





Effectiveness of round-robin depends on

- The number of jobs, and
- The size of the time quantum.

Large # of jobs means that the time between scheduling of a single job increases

(Slow responses)

Larger time quantum means that the time between the scheduling of a single job also increases (Slow responses)

Smaller time quantum means higher processing rates but also more overhead!

## Scheduling in general purpose systems

# **Priority scheduling**

Assign a priority (number) to each process. Schedule processes based on its priority.

Higher priority jobs processes get more CPU time.

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Managing priorities

Can use "nice" to reduce your priority

Can periodically adjust a process' priority

Prevents starvation of a lower priority process

Can improve performance of I/O-bound processes

p = aT_0 + (1-a)T_1

efficient to calculate when a = 1/2

Can use feedback to help determine the priorities

priority = 1/fraction last quantum used
```

## **Multi-Level Queue Scheduling**

Multiple queues, each with its own priority.

(Equivalently: Each priority has its own ready queue) Within each queue...

**Round-robin scheduling.** 

**Simplist Approach:** 

A Process's priority is fixed & unchanging



**Problem: Fixed priorities are too restrictive** 

Processes exhibit varying ratios of CPU to I/O times.

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Priorities are altered over time, as process behavior changes!



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**Dynamic Priorities** 

Priorities are altered over time, as process behavior changes!

**Issue:** When do you change the priority of a process and how often?

**Solution:** Let the amount of CPU used be an *indication* of how a process is to be handled

Expired time quantum  $\rightarrow$  more processing needed Unexpired time quantum  $\rightarrow$  less processing needed

Adjusting quantum and frequency vs. adjusting priority??

n priority levels, round-robin scheduling within a level
Quanta increase as priority decreases
Jobs are demoted to lower priorities if they do not complete within the current quantum



Details, details, details...

- Starting priority?
- High priority vs. low priority
- Moving between priorities?
- How long should the time quantum be?

# **Lottery Scheduling**

Scheduler gives each thread some lottery tickets.

To select the next process to run...

The scheduler randomly selects a lottery number

The winning process gets to run

*Example* Thread A gets 50 tickets

Thread B gets 15 tickets

Thread C gets 35 tickets

There are 100 tickets outstanding.

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**Thread B gets 15 tickets** 

Thread C gets 35 tickets  $\rightarrow$  35% of CPU

→ 15% of CPU

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There are 100 tickets outstanding.

- Flexible
- Fair
- Responsive.

Assume processes are relatively periodic Fixed amount of work per period (e.g. sensor systems or multimedia data)

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Two "main" types of schedulers...

**Rate-Monotonic Schedulers** 

**Earliest-Deadline-First Schedulers** 

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Two "main" types of schedulers...

**Rate-Monotonic Schedulers** 

Assign a *fixed, unchanging* priority to each process No dynamic adjustment of priorities

• Less aggressive allocation of processor

**Earliest-Deadline-First Schedulers** 

Assign dynamic priorities based upon deadlines.

Typically real-time systems involve several steps (that aren't in traditional systems)

Admission control

All processes must ask for resources ahead of time. If sufficient resources exist, the job is admitted "into" the system.

- Resource allocation
  - Upon admission...

the appropriate resources need to be reserved for the task.

Resource enforcement

Carry out the resource allocations properly.

For preemptable, periodic processes (tasks)

Assigns a fixed priority to each task

**T** = The period of the task

**C** = The amount of processing per task



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In RMS scheduling, the question to answer is...

What priority should be assigned to a given task?







The computation for P2 cannot be increased

Assign shortest period tasks to the highest priorities

**Admission control is difficult...** 

- < 70% There is a standard formula.
- > 70% It may be possible,

but need to do more complex analysis.

#### Assumptions:

- Processes complete (yield) within their period
- Independent processes
- Same CPU requirements per burst
- Other non-periodic processes have no deadlines

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• Instantaneous preemption with no overhead

When processes do not need periodic execution or have constant CPU requirements... When processes have deadline specifications...

#### Unlike RMS, EDF uses dynamic priorities (based upon earliest deadline first)

- (+) 100% processor utilization ...?
- (-) Need to keep track of deadlines

#### Admission Control

Just check to see if 100% processor utilization. Sum the  $C_i/T_i$ 's and see if less than or equal to 1 What about overhead?