

DPKernel

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	<i>TITLE :</i> DPKernel		
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REVISION HISTORY

NUMBER	DATE	DESCRIPTION	NAME

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Chapter 1

DPKernel

1.1 Library: DPKernel

SYSTEM DOCUMENTATION

Name: DPKERNEL
Version: 0.9 Beta
Date: December 1997
Author: Paul Manias
Copyright: DreamWorld Productions, 1996-1997. All rights reserved.
Notes: This document is still being written and will contain errors in a number of places. The information within cannot be treated as official until this autodoc reaches version 1.0.

CHANGES VERSION 0.9B

Added: Flush()
Read() (moved from files.guide)
Reset()
Write() (moved from files.guide)
Detach()
AutoStop()

Moved from the master.guide

FastRandom()
ReadKey()
SlowRandom()
WaitLMB()
WaitTime()
WriteDec()

Deleted: SetUserPrefs()
WaitLMB()
WriteDec()
ReadKey()
AutoSwitch()

Edited: WaitTime()
CloseDPK()
Activate()
AddSysObject()

1.2 DPKernel

OBJECTS

Module
RawData
Reference
Segment
SysObject
Task
Time

FUNCTIONS

AddSysEvent ()
AddSysObject ()
AllocMemBlock ()
AutoStop ()
CloseDPK ()
DPKForbid ()
DPKPermit ()
FastRandom ()
FindDPKTask ()
FindSysObject ()
FreeMemBlock ()
GetMemSize ()
GetMemType ()
OpenModule ()
RemSysEvent ()
SlowRandom ()
Switch ()
TagInit ()
TotalMem ()
WaitTime ()

Resource Tracking

AddTrack ()
DeleteTrack ()
InitDestruct ()
SelfDestruct ()

ACTIONS

Activate ()
Clear ()
CopyStructure ()
Deactivate ()
Detach ()
Display ()
Draw ()
Flush ()
Free ()
Get ()
Hide ()
Init ()
Load ()
Lock ()
Query ()
Read ()

```
Reset ()
Save ()
Unlock ()
Write ()
```

1.3 Kernel: Activate()

ACTION

```
Name:      Activate()
Short:     Perform the native action of the object.
Synopsis:  LONG Activate(APTR Object [a0]);
```

DESCRIPTION

Executes the native action for a given object (but only if the object supports Activate()). The action taken by each object class is extremely varied, here are some examples of what some current objects do for Activate():

```
Sound      - Plays sample data over the speakers.
Restore    - Replaces destroyed backgrounds.
Directory  - Grabs a list of directory contents.
```

Some objects will continue to perform their native action after they return back to you, for example a Sound will play in the background while you continue processing. If for some reason you wish to cancel an action, you can attempt a call to Deactivate().

INPUT

Object - Pointer to an initialised object structure.

RESULT

Returns ERR_OK on success.

SEE ALSO

Kernel: Deactivate()

1.4 Kernel: AddSysEvent()

FUNCTION

```
Name:      AddSysEvent ()
Short:     Adds a new event node to the system.
Synopsis:  *Event AddSysEvent(LONG *Tags)
```

DESCRIPTION

Note: This area of the OS is currently under development, so you use it at your own risk.

This function adds a new event node to the system. An event is a representation of a particular occurrence in the operating system. Such examples of this are the appearance of a new task, detecting a disk inserted by the user, low memory, a new screen on display, and so on.

Each event is handled by a particular area of the system. Sending floppy disk events is handled by the I/O routines for example. When a certain event is signalled by one of these handlers, the system will call all the nodes that are chained into that event. For an application or module to add a new node to one of these event chains, it needs to use `AddSysEvent()`.

Currently available events are:

```
OnNewTask
OnRemTask
```

EXAMPLE

This section of code will create a node that is linked to the `OnNewTask` event. Note that the node that is created will need to be removed with `RemSysEvent()` before the Task exits.

```
EVTOnNewTask = AddSysEventTags(TAGS, NULL,
    EVA_Number,    EVT_OnNewTask,    /* Event type */
    EVA_Routine,  &OnNewTask,       /* Pointer to event routine */
    EVA_Priority,  3,                /* Priority setting */
    TAGEND);
```

INPUT

Tags - Pointer to the event initialisation tags. See the Event object's documentation for details.

RESULT

Returns a pointer to the newly created Event node.

SEE ALSO

Kernel: `RemSysEvent()`

1.5 Kernel: AddSysObject()

FUNCTION

```
Name:      AddSysObject()
Short:     Adds a new object to the system.
Synopsis: *SysObject AddSysObject(APTR Tags [a0]);
```

DESCRIPTION

This function adds a new object to the system. The objective of this is to allow your object to support the various "Action" functions, and to give the system the necessary information about your object. Example:

```
JPEGObject = AddSysObjectTags(TAGS, NULL,
    SOA_ObjectID,    ID_HIDDEN,
    SOA_ClassID,    ID_PICTURE,
    SOA_Name,        "Picture~Jpeg",
    SOA_CopyToUniverse, &PIC_CopyToUniverse,
    SOA_CopyFromUniverse, &PIC_CopyFromUniverse,
    SOA_Load,        &PIC_Load,
    TAGEND);
```

In this example, `AddSysObject()` will create a Child type (called Jpeg) that belongs to the Picture class. Because child types inherit actions from the

parent, this Jpeg object will also support actions such as Query(), Init(), Get() and Free().

Master Objects

Any object that is initialised with a specific ID will be treated as a master object. If you call AddSysObject() when the given ID is already in use by another master object in the system, AddSysObject() will not create a new SysObject. Instead, it will add your object specifications to the existing SysObject.

For example, if you attempt to add a Picture master with a tag of SOA_Activate, and the existing Picture does not support that action, your supporting function will be added to the existing Picture master. If you specify SOA_Activate|TREPLACE, then your Activate function will be placed in the Picture master even if it already has support for Activate(). Please do not take the REPLACE flag lightly - only do this if absolutely necessary.

Hidden Objects

You can declare a "hidden" system object by passing ID_HIDDEN. What this creates is an object that has no identification characters. This means that no program will be able to specifically search for your object in the system. For example, the Get() action will be completely useless in looking for your object. The only way to get to your object will be via the FindSysObject() function.

The main purpose for creating hidden objects is for simple object types that will only be used internally (in the system). For example, XPK support was implemented via a hidden object, because as a programmer you wouldn't want to deal with unpacking when loading data. Instead, the Load() function will invisibly detect XPK files for you and call upon the hidden XPK object to deal with the data.

Child Objects

To declare a child object, use ID_HIDDEN in SOA_ObjectID, then use the class identifier in SOA_ClassID (e.g. ID_PICTURE). This is useful in situations where you might want to add support for a new data format to an object. As an example, when a jpeg picture is loaded the master will fail to recognise it, because it only uses IFF files. This failure will cause the Load() function to look for child objects belonging to the Picture class and use those instead.

The other advantage as a child object is that you get function inheritance from the master. So if any action is unsupported by your object, you will still be able to inherit support from the master's functions.

INPUT

Tags - Pointer to a standard tag list.

RESULT

Returns a pointer to the created SysObject or NULL if failure.

SEE ALSO

Kernel: Free()

1.6 Kernel: AddTrack()

FUNCTION

Name: AddTrack()

Short: Adds a resource tracking node.

Synopsis: LONG AddTrack(LONG Resource [d0], LONG Data [d3], void *Routine [a0])

DESCRIPTION

This function is intended for use by the system modules, but you may use it in your program if necessary.

About Resource Tracking

Each task has a special list of all system resources that are currently in use. Each resource has its own list node that uniquely identifies it (eg each memory allocation has its own resource node). When a particular resource has been freed the list is checked and its related node is deleted. When the program shuts down, any resources that are still in the list can be freed by the kernel so that they are not lingering in the system.

When you call AddTrack() you need to give it a resource ID which is made up from one of the following resource types. This is important as resources are freed from the system in an order based on these ID's. The order looks like this:

1. Free all hardware based resources (blitter, sound, etc).
2. Free complex resources (both hardware and software).
3. Free customised resources (user defined types etc).
4. Free memory (always freed last).

As you can see the correct order is vital, if memory was freed first it would have adverse affects when freeing the complex and user resource types. Always make sure that you give the correct ID when describing your resource. The ID's are RES_MEMORY, RES_HARDWARE, RES_COMPLEX and RES_CUSTOM, as outlined in the include file "system/tasks.i".

Passing a Data pointer is optional, but you will need it to uniquely identify your resource later on (eg AllocMemBlock() uses this to store a pointer to the allocated memory).

The Routine points to code that will free the resource if it is still in use when the program exits. It will be passed the Data field of the resource node in register d0, and the Key field in register d1. This routine must save all the CPU registers that it uses and can only free its resource - it may not make new allocations of any sort. It may send debug and error messages, which will help to work out what the problem is if anything goes wrong.

INPUT

Resource - Correct resource identifier from system/tasks.i.

Data - Optional data pointer to store in the resource node.

Routine - Pointer to the routine to call when freeing the resource.

RESULT

Returns a key that AddTrack() has used to uniquely identify the resource. You will need to store the key somewhere for when you call DeleteTrack().

SEE ALSO

Kernel: DeleteTrack()

Include: system/tasks.i

1.7 Kernel: AllocMemBlock()

FUNCTION

Name: AllocMemBlock()

Short: Allocate a new memory block.

Synopsis: APTR AllocMemBlock(LONG Size [d0], LONG MemType [d1]);

DESCRIPTION

This function allocates a memory block from the system memory pool and returns it to your program. By default all memory is public. This open memory model is used to enhance the communication levels between tasks and functions. If you need private memory make sure that you ask for it and that you will be the only Task using it. An AllocPrivate() macro has been included to make this type of allocation easier.

Protection and Resource Tracking

Header and Tail ID's are used to offer a security system similar to MungWall, acting as cookies at each end of a memory block. You will be alerted by FreeMemBlock() if the ID's are damaged. This is a permanent debugging feature, so there is little need to run MungWall for debugging your programs.

Resource tracking is automatic, so you will be warned if you allocate memory and forget to free it on exit (ie when you close down). Any memory that is found will be freed for your convenience.

By default all memory is cleared before it is given to you. Here are the memory types:

MEM_DATA

Suitable for storing data and variables. This is the default. Note that you are disallowed from running code from this type of memory.

MEM_CODE

This can store and execute CPU instructions. It also fits the requirements of MEM_DATA, so you can store variables and data in it.

MEM_VIDEO

Is for displaying graphics, and is also compatible with the Blitter interface.

MEM_BLIT

Is memory that is compatible with the blitter module. Currently this module only uses chip memory, but future versions could also support CPU drawing from fast if the graphic is located in that area.

MEM_SOUND

For memory that is compatible with the Sound interface.

You may also use the following extra flags when making your memory

allocation:

MEM_PRIVATE

If other programs should not have access to your memory.

MEM_UNTRACKED

If you don't want resource tracking on your allocation. You should only use this flag if the memory is part of a complex resource that is already being tracked.

INPUT

Size - Size of the required memblock in bytes.

MemType - The type of memory to allocate, eg MEM_VIDEO.

RESULT

Pointer to the start of your allocated memblock or NULL if failure.

SEE ALSO

Kernel: FreeMemBlock()

GetMemSize()

GetMemType()

1.8 Kernel: AutoStop()

FUNCTION

Name: AutoStop()

Short: Pauses your task when the user leaves your screen or windows.

Synopsis: void AutoStop(void)

DESCRIPTION

If the user moves away from your task by switching to a different screen or window, you may want to pause your actions until the user returns the focus back to your task. This is very useful in games and other "non-stop" programs which can continue regardless of any user interaction. Such non-stop programs can be counter productive if the user is at a different task, and your program is stealing resources in the background.

If this function finds that the user has moved away, it pauses your task immediately. When the user returns to your task, this function returns to execute the next instruction in your program.

SEE ALSO

Kernel: Switch()

1.9 Kernel: Clear()

ACTION

Name: Clear()

Short: Clears an object's graphic from its container.

Synopsis: LONG Clear(APTR Object [a0])

DESCRIPTION

This action will clear an object's graphic from its container. For example, `Clear(Bob)` will clear a Bob graphic from its assigned Bitmap. Note that it will not restore what was previously under the graphic - it just clears it all away. This means that you will get a empty rectangle or mask of the object, most likely with a colour of black.

INPUT

Object - Pointer to an initialised object.

RESULT

Returns `ERR_OK` if successful.

SEE ALSO

Kernel: `Draw()`

1.10 Kernel: `CloseDPK()`

FUNCTION

Name: `CloseDPK()`
Short: Closes the kernel library.
Synopsis: `void CloseDPK(void)`

DESCRIPTION

Before your program exits you will have to call this function so that the system knows you are shutting down. If you do not close the kernel before you exit you will leave certain memory allocations unfreed and there may be other adverse system effects.

This function will perform a resource tracking check, so if you have not freed any system resources you will be notified here (if you get a yellow alert box, you will need to use `IceBreaker` to get a detailed list of the errors).

`CloseDPK()` is responsible for handling the `OnRemTask` event. This event will be activated as soon as `CloseDPK()` is called by your task.

NOTE

Remember that you may not call any more functions after calling `CloseDPK()`.

This function is used in the `STARTDPK` macro and the `dpk.o` file, so C and assembler programmers do not need to call this function explicitly. Programs that are started natively or from the `StartDPK` executable, will not have to call `CloseDPK()`. This is because the call will be handled within the system for these cases.

1.11 Kernel: `CopyStructure()`

ACTION

Name: `CopyStructure()`
Short: Copies details from one structure to another.
Synopsis: `LONG CopyStructure(APTR Source [a0], APTR Destination [a1]);`

DESCRIPTION

Copies one structure's data across to another. The structures can be similar or completely different (even a Screen -> Sound copy is possible, although there won't be much of a result).

This action only copies fields that are currently present in the Universe object. This action DOES NOT copy fields that contain option flags, fields that contain pointers to the main area of data (eg screen bitplanes), or fields designated as private.

Only the NULL fields in the Destination structure will be written to. If the Destination structure has already been initialised, you may find that CopyStructure() has no effect due to this condition.

INPUT

Source - Points to the source object.
Destination - Points to the destination object.

RESULT

Returns ERR_OK on success, otherwise may fail due to unrecognised objects or invalid arguments.

1.12 Kernel: Deactivate()

FUNCTION

Name: Deactivate()
Short: Stop the object from performing its native action.
Synopsis: void Deactivate(APTR Object [a0]);

DESCRIPTION

Stops the given object from continuing its native action. This function only works if the object has recently had Activate() called on it. If the object is not performing any actions, then this function will simply do nothing.

INPUT

Object - Pointer to an initialised object structure.

SEE ALSO

Kernel: Activate()

1.13 Kernel: DeleteTrack()

FUNCTION

Name: DeleteTrack()
Short: Delete a resource tracking node.
Synopsis: void DeleteTrack(LONG Key [d1]);

DESCRIPTION

Deletes a resource node allocated from AddTrack(). If resource nodes are not deleted they will stay linked to the program and the system will attempt to free them when the program shuts down.

Note that this function only deletes the resource node, it will not attempt to deallocate the resource from its deallocation function.

INPUT

Key - A key that was obtained from AddTrack().

SEE ALSO

Kernel: AddTrack()

1.14 Kernel: Detach()

ACTION

Name: Detach()

Short: Detach an object from a parent structure.

Synopsis: LONG Detach(APTR Object [a0], APTR Parent [a1])

DESCRIPTION

This action will separate an object from a parent structure, without causing any undue harm to either object. If this operation is successful, you can then use the detached object independantly of its parent.

Some detachments are impossible or too dangerous to attempt, e.g. a Screen object cannot survive without its Bitmap.

INPUTS

Object - The object that needs to be detached.

Parent - Pointer to the parent object that you want to be detached from.

RESULT

Returns ERR_OK if successful. This call can fail if the object does not support the Detach() action, or if it is not attached to a parent object in the first place.

1.15 Kernel: Display()

ACTION

Name: Display()

Short: Displays an object inside its Container.

Synopsis: LONG Display(APTR Object [a0]);

DESCRIPTION

Calling this action will display the object in its specified container. If the object cannot be displayed, or if the required container was never specified on Init(), the the call will fail.

INPUT

Object - The object to be displayed.

RESULT

Returns ERR_OK on success.

SEE ALSO

Kernel: Hide()

1.16 Kernel: DPKForbid()

FUNCTION

Name: DPKForbid()

Short: Stop other tasks/processes from executing.

Synopsis: void DPKForbid(void)

DESCRIPTION

Stops all other tasks and processes from executing until you call DPKPermit(). This call will not turn off interrupts.

NOTE

This function has little effect in systems that do not multi-task.

SEE ALSO

Kernel: DPKPermit()

1.17 Kernel: DPKPermit()

FUNCTION

Name: DPKPermit()

Short: Allow other tasks to continue their processing.

Synopsis: void DPKPermit(void);

DESCRIPTION

Reverses a previous call to DPKForbid(), so that all tasks can continue their normal processes.

SEE ALSO

Kernel: DPKForbid

1.18 Kernel: Draw()

ACTION

Name: Draw()

Short: Draws an object's graphic to its container.

Synopsis: LONG Draw(APTR Object [a0]);

DESCRIPTION

This action will draw an object's graphic to the container that it was originally initialised to. If the object was never initialised to a container that supports drawing, this action will have no effect. Secondly, if the object has no graphical representation, then nothing can be drawn.

This action is most often used in the drawing of Bobs and MBobs.

INPUT

Object - Points to an initialised object that you wish to draw.

RESULT

Returns ERR_OK if successful.

SEE ALSO

Kernel: Clear()

1.19 Kernel: FastRandom()

FUNCTION

Name: FastRandom()

Short: Generate a random number between 0 and <Range>.

Synopsis: LONG FastRandom(WORD Range [d1]);

DESCRIPTION

Creates a random number as quickly as possible. The routine uses one divide to determine the range and will automatically change the random seed value each time you call it. This routine has now been fully tested and generates 100% patternless numbers.

Remember that all generated numbers fall BELOW the Range. Add 1 to your range if you want this number included.

INPUTS

Range - A range between 1 and 32767. An invalid range of 0 will result in a division by zero error.

RESULT

A number greater or equal to 0, and less than Range.

SEE ALSO

Kernel: SlowRandom()

Demos: demos/randomplot

1.20 Kernel: FindDPKTask()

FUNCTION

Name: FindDPKTask()

Short: Find the DPKTask structure for the current task.

Synopsis: *DPKTask FindDPKTask(void)

DESCRIPTION

This function will return the DPKTask structure for the task that called it. The DPKTask structure is used for storing data that is specific to your task - things like preference settings for example. Almost all of the DPKTask fields are private and you cannot write to this structure unless you are a system module.

For the curious, it only takes 3 assembler instructions to grab the task node, so there is no time wasted in calling this function.

RESULT

Pointer to the DPKTask structure.

SEE ALSO

Include: system/tasks.i

1.21 Kernel: FindSysObject()

FUNCTION

Name: FindSysObject()

Short: Finds a system object based on the ID.

Synopsis: *SysObject FindSysObject(LONG ID [d0], *SysObject [a0]);

DESCRIPTION

This function begins finding a SysObject structure when given an ID and an initial SysObject of NULL. If a matching SysObject is found in the system, it will be returned immediately. Otherwise this function will go through the object reference list to see if the object is available on disk. If this is the case, then the relevant module will be loaded and then the SysObject will be returned. If not, NULL is returned to indicate a failure in the search.

The first matching SysObject to be returned is always the master object of that particular class. If you want to find the child objects of the class, call FindSysObject() again and supply the previously returned SysObject. The next matching object will be found and returned, otherwise NULL is returned if no more objects are left. Note that if you are looking for a specific child or hidden object you will need to check the SysObject->Name string.

NOTE

The only way to find hidden objects are to search on ID_HIDDEN.

INPUT

ID - A system object ID as specified in system/register.i.

SysObject - Last received SysObject if continuing a search, otherwise NULL to start a new search.

RESULT

Pointer to the SysObject that has been found as a result of the search, or NULL if no matching objects were found.

SEE ALSO

Include: system/register.i

1.22 Kernel: Flush()

ACTION

Name: Flush()

Short: Flush buffered data from an object.

Synopsis: LONG Flush(APTR Object [a0])

DESCRIPTION

This action will flush all buffered data from an object. It is mostly intended for file objects that may buffer data to speed up processing time. As a result of flushing, any un-written data will be dumped to its physical location, if this is practical for the object in question (eg files).

Note that any object using a buffering technique will automatically flush its data when you Free() it.

INPUT

Object - Pointer to an initialised object.

RESULT

Returns ERR_OK on success. All errors returned from Flush() are non-fatal, so you may continue to use the object after failure.

SEE ALSO

Kernel: Reset()

1.23 Kernel: Free()

ACTION

Name: Free()

Short: Frees an object and any of its associated parts.

Synopsis: void Free(APTR [a0])

DESCRIPTION

This action will take any system object and free its resources. It accepts lists for multiple deallocations.

If there is no Free action for the object, then it will be assumed that the object does not require freeing. However, such an instance would be highly unusual.

Once the object has been freed, the original pointer that you passed to Free() immediately becomes invalid. For this reason, any references to the object should be driven to NULL to prevent bugs from appearing in your program.

INPUT

APTR - Pointer to one of the following:

Object, ListV1, ListV2, TagList, ObjectList

SEE ALSO

Kernel: Get()

Init()

1.24 Kernel: FreeMemBlock()

FUNCTION

Name: FreeMemBlock()
Short: Free a previously allocated mem block.
Synopsis: void FreeMemBlock(APTR MemBlock [d0])

DESCRIPTION

Frees a memory area allocated by AllocMemBlock(), AllocVideoMem(), AllocBlitMem(), or AllocSoundMem(). If the mem header or tail is missing, then it is assumed that something has written over the boundaries of your memblock, or you are attempting to free a non-existent allocation. Normally this would cause a complete system crash, but instead we simply send a message to IceBreaker and leave the memory block in the system.

Bear in mind that it does pay to save your work and reset your machine if such a message appears, as it indicates that important memory data may have been destroyed.

NOTE

Never attempt to free the same MemBlock twice.

INPUT

MemBlock - Points to the start of a memblock. If NULL, then no action will be taken (function exits).

SEE ALSO

Kernel: AllocMemBlock()

1.25 Kernel: Get()

FUNCTION

Name: Get()
Short: Gets the latest version of a specified object.
Synopsis: APTR Get(LONG ID [d0])

DESCRIPTION

This function will get the latest version of any object that you specify by the ID argument. That is of course, if the object has been correctly installed and is registered within the systsem.

Screens, Pictures and Sounds are permanent and therefore always recognised, while something like a CardSet is an extra and must have been installed first [Get() will find such objects by using the reference files in GMS:System/References/].

All objects going through Get() are tagged with a resource key. Before your program exits you will need to free the object with the Free() action. If you want your program to exit and leave certain objects in the system (ie for other programs to use), then you can logical OR the ID argument with GET_NOTRACK.

The structure will return with empty fields, so you can fill them out to suit your requirements.

This function is the only way you can legally obtain a system object. If

you ever think of compiling objects directly inside your program, forget it
- your program will crash.

INPUT

ID - One of the ID's as specified in the system/register.i file.

RESULT

The latest version of the specified object or NULL if failure (caused by lack of memory or unrecognised ID).

SEE ALSO

Kernel: Free()

Init()

Include: system/register.i

1.26 Kernel: GetMemSize()

FUNCTION

Name: GetMemSize()

Short: Identifies the size of a given memory block.

Synopsis: LONG GetMemSize(APTR MemBlock [a0])

DESCRIPTION

This function will get the size of any memory block legally obtained from AllocMemBlock(). Illegal pointers will result in a return of NULL.

INPUT

MemBlock - Pointer to the start of the memory block to be identified.

RESULT

The size of the memory block in bytes.

SEE ALSO

Kernel: GetMemType()

1.27 Kernel: GetMemType()

FUNCTION

Name: GetMemType()

Short: Identifies the type of memory in use by a particular memory block.

Synopsis: LONG GetMemType(APTR MemBlock [a0])

DESCRIPTION

This function will get the memory type of any memory block legally obtained from AllocMemBlock(). Illegal pointers will result in a return of -1.

INPUT

MemBlock - Pointer to the start of the memory block to be identified.

RESULT

Type - Memory flags identifying the block (see AllocMemBlock()).

SEE ALSO

Kernel: GetMemSize()

1.28 Kernel: Hide()

ACTION

Name: Hide()

Short: Hides a displayed object from view.

Synopsis: void Hide(APTR Object [a0])

DESCRIPTION

If you have successfully displayed an object, then you can call Hide() to take it off the display. If the object has not been displayed, then the call is ignored.

INPUT

Object - Pointer to an initialised object that has been displayed.

SEE ALSO

Kernel: Display()

1.29 Kernel: Init()

ACTION

Name: Init(), InitTags()

Short: Initialises an object so that it is ready for active use.

Synopsis: APTR Init(APTR Data [a0], APTR Container [a1])
APTR InitTags(APTR Container, LONG tagltype, ...)

DESCRIPTION

This function initialises any recognised system object. The container argument is dependent on the type of object that you are initialising (eg Bob requires a Screen or Bitmap) so it is not always necessary to supply one.

If you provide a List or ObjectList then be aware that Init() will pass the original container to ALL the objects in the list. For this reason all objects must share some commonalities to the container (eg do not initialise a Sound in a list of Bobs).

If the initialisation action fails, then Init() will look for any child classes that can handle initialisation of the object. If a child class succeeds it will gain full ownership of the object and the master will lose it.

This action will return NULL on error, you can get very informative error messages by using IceBreaker.

Note to Module Programmers

If your module fails to initialise the object, Init() will free the object for you - do not attempt to free the object yourself.

INPUT

Data - Pointer to an: Object, TagList, ListV1, ListV2 or an ObjectList.
Container - Some objects need to be initialised to a "container" or parent object. If this is the case, specify that object here.

RESULT

Pointer to the initialised object.

SEE ALSO

Kernel: Free()
Get()

1.30 Kernel: InitDestruct()

FUNCTION

Name: InitDestruct()
Short: Initialise the task for use of SelfDestruct().
Synopsis: void InitDestruct(APTR DestructCode [a0], APTR DestructStack [a1])

DESCRIPTION

This is a special function that is called in the STARTDPK macro, gms.o startup file and StartDPK program only. You should never call this function explicitly unless you are writing your own startup code.

InitDestruct() will prepare your task so that it may be destroyed by the SelfDestruct() function. DestructCode must point to the exit code for your task. The exit code must call CloseDPK() at some point if you are to free your tasks resources. DestructStack must point to the correct stack area for your exit code, otherwise your task cannot return to the system correctly.

To see an example of how this function works look at the STARTDPK macro in file "dpkernel.i".

INPUTS

DestructCode - Points to the code at which your task makes its exit.
DestructStack - Points to the stack that will be used for the exit.

SEE ALSO

Kernel: SelfDestruct()

1.31 Kernel: Load()

ACTION

Name: Load()
Short: Load a file and initialise it as a system object.
Synopsis: APTR Load(APTR Source [a0], LONG ID [d0])

DESCRIPTION

Loads in a file, finds its native object and then returns an object that has been initialised and ready for use.

If you supply a special ID of ID_MEMBLOCK, Load() will allocate a MEM_DATA memory block, load all of the files contents into it, and then return the memory block pointer back to you. You will need to free this pointer with FreeMemBlock() when you are finished with it.

The file will be loaded in according to the object's preferred format, eg loading of Pictures will always result in the data being in video ram. If you require more power in the loading of a particular object, use the Init() function instead of Load().

NOTE

If this function cannot find an object that immediately recognises this structure, Load() will most probably open the file using the RawData object.

INPUT

Source - File name or Memory location pointer.
ID - Forces the type of object that you want to be returned. Use NULL if you want Load() to assess the file and return its appropriate object.

RESULT

Pointer to the initialised object or NULL if failure. Remember that RawData is used if no other object recognises the source file.

SEE ALSO

Kernel: Free()
Init()

1.32 Kernel: OpenModule()

FUNCTION

Name: OpenModule()
Short: Provides a quick way of opening a module.
Synopsis: *Module OpenModule(LONG ID [d0], BYTE *Name [a0])

DESCRIPTION

INPUT

ID - ID of the module to open (can be NULL if name is supplied).
Name - Name of the module to open (not required if ID is known).

RESULT

Pointer to an initialised Module structure or NULL if failure.

SEE ALSO

Object: Module
Kernel: Init()

1.33 Kernel: Query()

ACTION

Name: Query()

Short: Gets the latest information on a particular object.

Synopsis: LONG Query(APTR Object [a0])

DESCRIPTION

Calling the Query() action will update all fields in the given object to reflect any changes since the last Query() or initialisation.

Examples of using the Query() action are to get information on picture files (such as width, height, amount of colours) and obtaining consistently changing data, such as joystick information.

INPUT

Object - Pointer to an object allocated from Get().

RESULT

ErrorCode - ERR_OK on success.

1.34 Kernel: Read()

ACTION

Name: Read()

Short: Read data from an object into a buffer.

Synopsis: LONG Read(APTR Object [a0], APTR Buffer [a1], LONG Length [d0]);

DESCRIPTION

This action will read the amount of bytes as determined by Length, from the Object and into the given memory Buffer. The read will start at the position determined by the object's Byte Position field. This field will be incremented to BytePos+Length if the call succeeds. Further calls to Read() will therefore start from where you left off.

If the Length exceeds the total size of the data then this function will only read as many bytes as there are left in the file. You can always compare the BytePos and Size fields to see how many bytes are left to read in the object.

INPUT

Object - Pointer to an initialised Object.

Buffer - Pointer to a memory area in which the data will be written to.

Length - Amount of bytes to read from the object.

RESULT

Returns the total amount of data read into the buffer. NULL can indicate an error, or maybe there is no more data left to read.

SEE ALSO

Kernel: Write()

1.35 Kernel: RemSysEvent()

FUNCTION

Name: RemSysEvent()
Short: Removes an event tag from the system.
Synopsis: void RemSysEvent(*Event [a0])

DESCRIPTION

Not documented yet.

INPUT

Event -

SEE ALSO

Kernel: AddSysEvent()

1.36 Kernel: Reset()

ACTION

Name: Reset()
Short: Resets an object to a receptive state.
Synopsis: LONG Reset(APTR Object [a0])

DESCRIPTION

This action will reset an object to a state that is receptive to new processes. Its behaviour depends greatly on the object at hand. Files will reset their byte position to 0 for all further I/O operations, while the Restore object dumps all current restore states and forgets about any buffered data.

NOTE

Reset() is not the same as Flush(). Some objects will flush themselves before performing the reset action, others may not. It depends on the circumstances, but if you want the data flushed then use the Flush() action.

INPUT

Object - Pointer to an initialised object.

RESULT

An error will be returned if the action could not execute. This can happen if the object does not support or require the Reset() action.

SEE ALSO

Kernel: Flush()

1.37 Kernel: Save()

ACTION

Name: Save()
Short: Writes all of an object's information to a destination file.
Synopsis: LONG Save(APTR Object [a0], APTR Destination [a1])

DESCRIPTION

This action will save an object's data to a file, in a data format that is suitable for the Load() action. For example: As a default, the Picture object saves all pictures in IFF format, and it also supports the loading of pictures in IFF format.

It will be possible for a child object to save a file in its own format (ie instead of the master format). This way the user can perform file conversions, such as load a picture as IFF and save it out as JPEG. The method for doing this is still under-way, but to simplify the process for programmers, the user will probably be able to choose preferred saving types from GMSPrefs.

INPUT

Object - Pointer to an initialised object.
Destination - Pointer to a Source/Destination structure, such as a FileName, MemPtr or File object.

RESULT

Returns ERR_OK on success.

SEE ALSO

Kernel: Load()

1.38 Kernel: SelfDestruct()

FUNCTION

Name: SelfDestruct()
Short: Destroys the task and frees resources.
Synopsis: void SelfDestruct(void)

DESCRIPTION

Destroys the task that called this function and then proceeds to free all of its resources according to the resource nodes. This is a completely safe and effective way of destroying a task, and can be used for deconstructing a task when it has got into unrecoverable circumstances.

You must have called InitDestruct() before calling this function. If you are programming in C or assembler this initialisation is already in the STARTDPK macro and gms.o startup file, so this does not concern you.

NOTE

This function will not return. However if InitDestruct() has not been called then the function will not be able to do anything and will return back to the task.

SEE ALSO

Kernel: InitDestruct()

1.39 Kernel: SlowRandom()

FUNCTION

Name: SlowRandom()
Short: Generate a random number between 0 and <Range>.
Synopsis: LONG SlowRandom(WORD Range [d1])

DESCRIPTION

Generates a very good random number in a relatively short amount of time. This routine takes approximately two times longer than FastRandom(), but is guaranteed of giving excellent random number sequences.

Remember that all generated numbers fall BELOW the Range. Add 1 to your range if you want this number included.

INPUTS

Range - A range between 1 and 32767.

RESULT

A number greater or equal to 0, and less than Range.

SEE ALSO

Kernel: FastRandom()
Demos: demos/randomplot

1.40 Kernel: Switch()

FUNCTION

Name: Switch()
Short: Stops the task that called this function.
Synopsis: void Switch(void);

DESCRIPTION

Switches your task over to the next task in the queue. This function will not return until the user reactivates your task, so your task's execution is effectively stopped. Any secondary processes and interrupts that you have spawned will continue to execute, so multi-tasking can still be effective.

If the next task is screen-based, then your screen display will be removed and the new screen will be displayed. If you have any secondary tasks running, then take note: You must not allow them to use the drawing/blitter operations as your display memory may be temporarily moved to free up video memory. Blitting to an invisible display is also considered to be bad practice as most GMS tasks require all available blitter time. We also ask you to refrain from using the audio functions as the next task will probably be needing all available channels.

If there are no more tasks in the queue, then the screen display will return to intuition. GMS supports two methods of screen switching to intuition, Switch-To-Window and Switch-To-Screen. The method used depends on the setting in the GMSPrefs utility.

Switch-To-Window drops out to workbench and places a window on the screen. It will wait until the close gadget is pressed, whereupon your game will continue where it left off.

Switch-To-Screen opens an intuition screen and busy-waits until that screen comes to the front. At that point the intuition screen will be closed and your game will resume execution.

SEE ALSO

Kernel: AutoStop()

1.41 Kernel: TagInit()

FUNCTION

Name: TagInit()

Short: Initialise a structure according to a tag list.

Synopsis: LONG TagInit(APTR Structure [a0], APTR TagList [a1])

DESCRIPTION

This function is intended for system modules but may be used by normal programs if required.

It will process a standard tag list and store specified values in the given structure, which should be empty although this is not a pre-requisite. It is important that the tags themselves have been correctly defined using the TBYTE, TWORD and TLONG flags. Check the include files for examples.

This function has some software based memory protection and will prevent values from being written outside of the structure's memory area. Detected errors will be sent to the system debugger.

INPUTS

Structure - Pointer to allocated structure memory.

TagList - Pointer to a standard tag list (see tags).

RESULT

Returns ERR_OK if successful.

SEE ALSO

Kernel: Get()

1.42 Kernel: TotalMem()

FUNCTION

Name: TotalMem()

Short: Gets the total amount of memory used by a task.

Synopsis: LONG TotalMem(*DPKTask [a0], LONG Flags [d0])

DESCRIPTION

Gets the total amount of memory currently in use by a Task. This total can be calculated from a specific type of memory, or a total of all memory in use if -1 is specified. If DPKTask is passed as NULL then this function will sum up the total amount of memory used by all Tasks in the system.

INPUT

DPKTask - Pointer to a DPKTask object, or NULL to calculate from all tasks.

Flags - The memory type that you want to be calculated (MEM_VIDEO, MEM_SOUND, MEM_DATA or MEM_BLIT). If you want a complete total of all memory in use, specify -1 here.

RESULT

The total amount of memory in use.

1.43 Kernel: WaitTime()

FUNCTION

Name: WaitTime()

Short: Wait for a specified amount of micro-seconds.

Synopsis: void WaitTime(LONG MicroSeconds [d0])

DESCRIPTION

Waits for a specified amount of micro-seconds. During this time it will reduce the task priority and make regular calls to AutoStop() for you.

INPUT

MicroSeconds - Amount of micro-seconds to wait for (100 = 1 Second).

1.44 Kernel: Write()

ACTION

Name: Write()

Short: Writes a given amount of bytes to an object's data space.

Synopsis: LONG Write(APTR Object [a0], APTR Buffer [a1], LONG Length [d0])

DESCRIPTION

This action will write the amount of bytes as determined by 'Length', from the memory buffer and into the Object's data space. The write will start at the position determined by the object's BytePos field. This field will be incremented to BytePos+Length if the call succeeds.

If the BytePos+Length exceeds the total size of the file, then the object may increase its data space to cope with writing out the rest of the buffer. This will always happen with Files, but memory based objects will rarely be able to do this.

INPUTS

Object - Pointer to an initialised object.

Buffer - Pointer to data that will be written to the object.

Length - Amount of bytes to write.

RESULT

Returns the amount of bytes written. NULL indicates an error.

SEE ALSO

Kernel: Read()

1.45 Kernel: ()

FUNCTION

Name: ()

Short:

Synopsis:

DESCRIPTION

NOTE

INPUT

RESULT

SEE ALSO
