

How Messy is Your Database?













- ◆ **A Robelle and Adager Tutorial**
HP World '96
Anaheim, California
August 5 - 8, 1996



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How messy is your database?

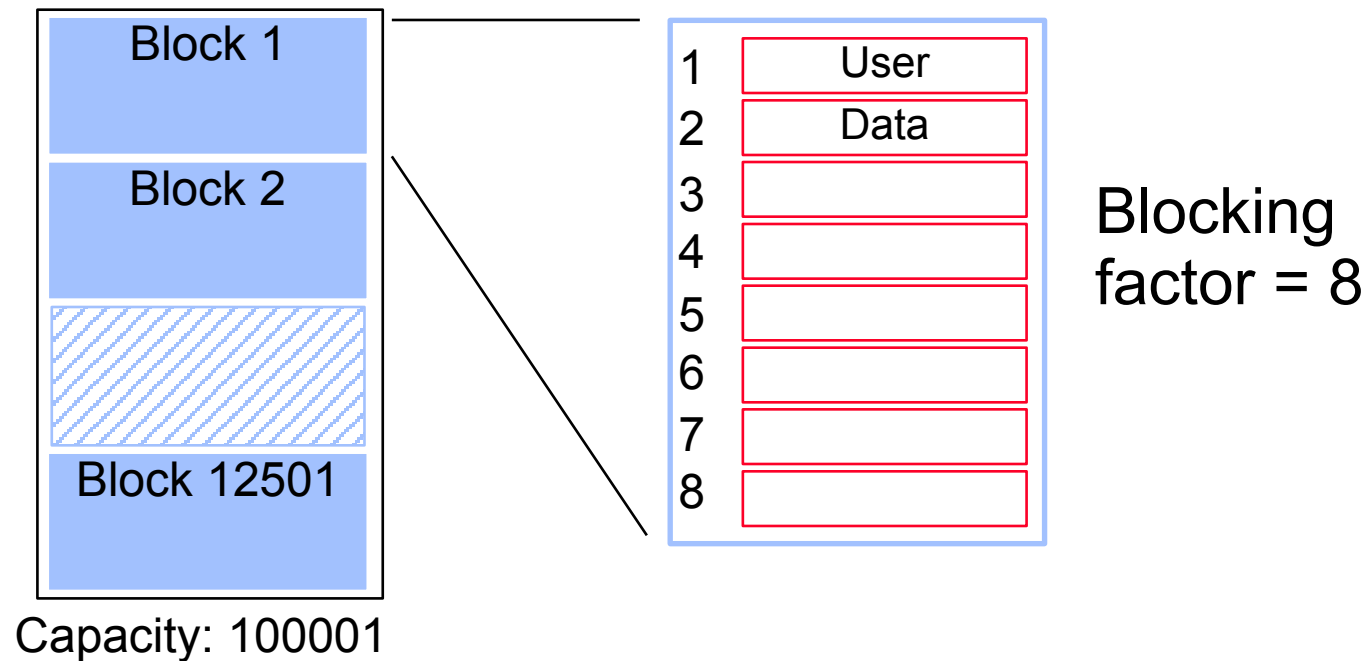


- A database is messy if it takes more I/O than it should
- Unnecessary I/O is still a major limiting factor even on MPE/iX machines
- Databases are messy by nature
- Run HowMessy or DBLOADNG against your database
 - HowMessy is a bonus program for Robelle customers
 - DBLOADNG is a contributed library program

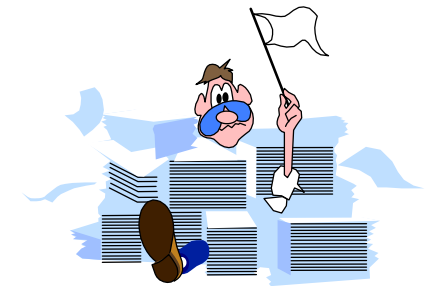
Blocks



- ☐ TurboIMAGE does all I/O operations in blocks
- ☐ A block may contain many user records
- ☐ More entries per block means fewer I/Os
- ☐ Fewer I/Os means better performance

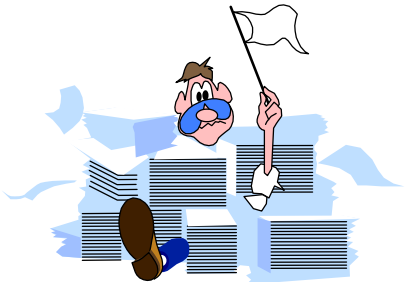


Record location in masters

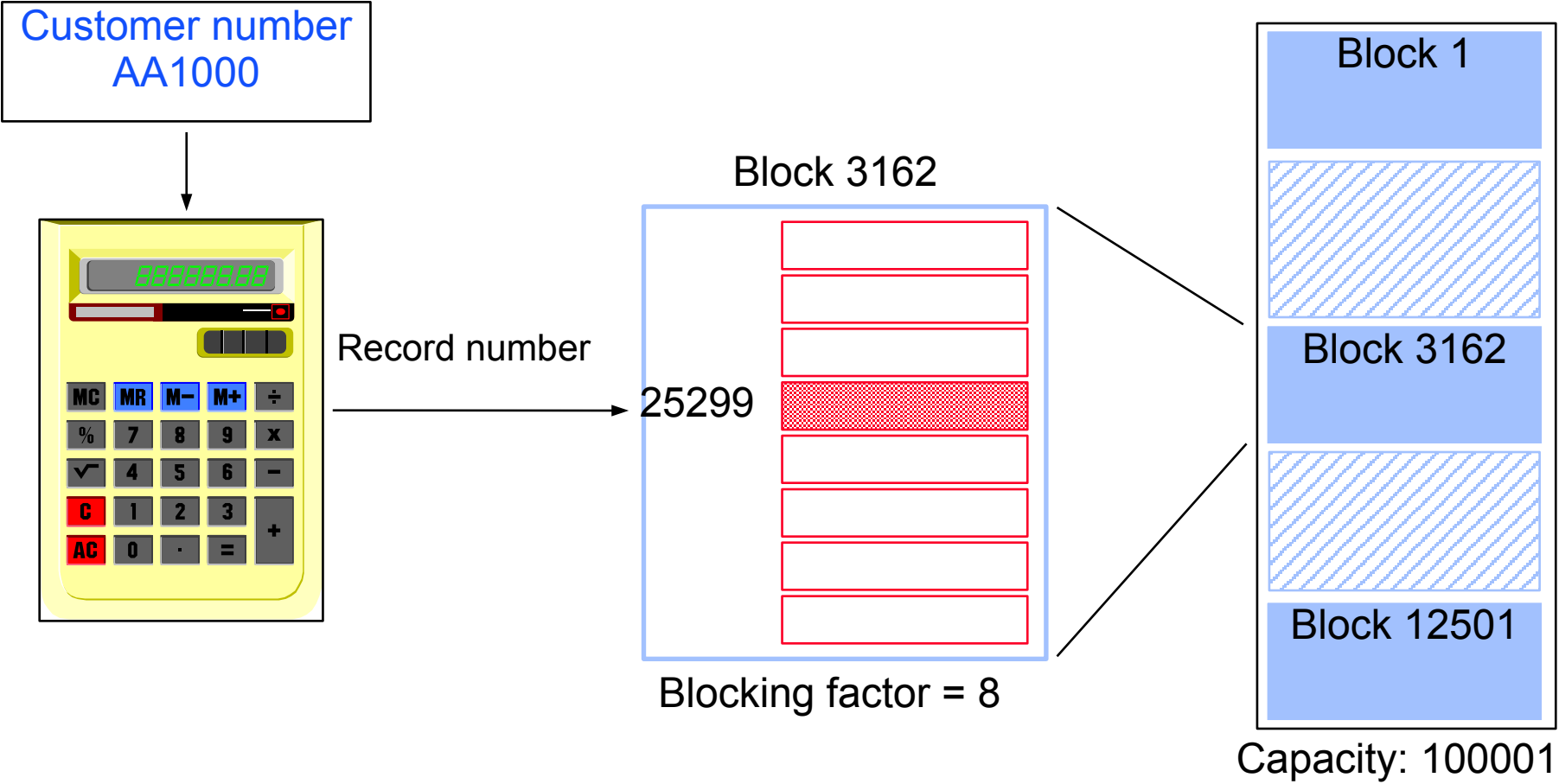


- ❑ Search item values must be unique
- ❑ Location of entries is determined by a hashing algorithm or a primary address calculation
- ❑ Calculation is done on search item value to transform it into a record number between one and the capacity
- ❑ Different calculation depending on the search item type
 - ❑ X, U, Z, and P give random results
 - ❑ I, J, K, R, and E give predictable results

Hashing algorithm



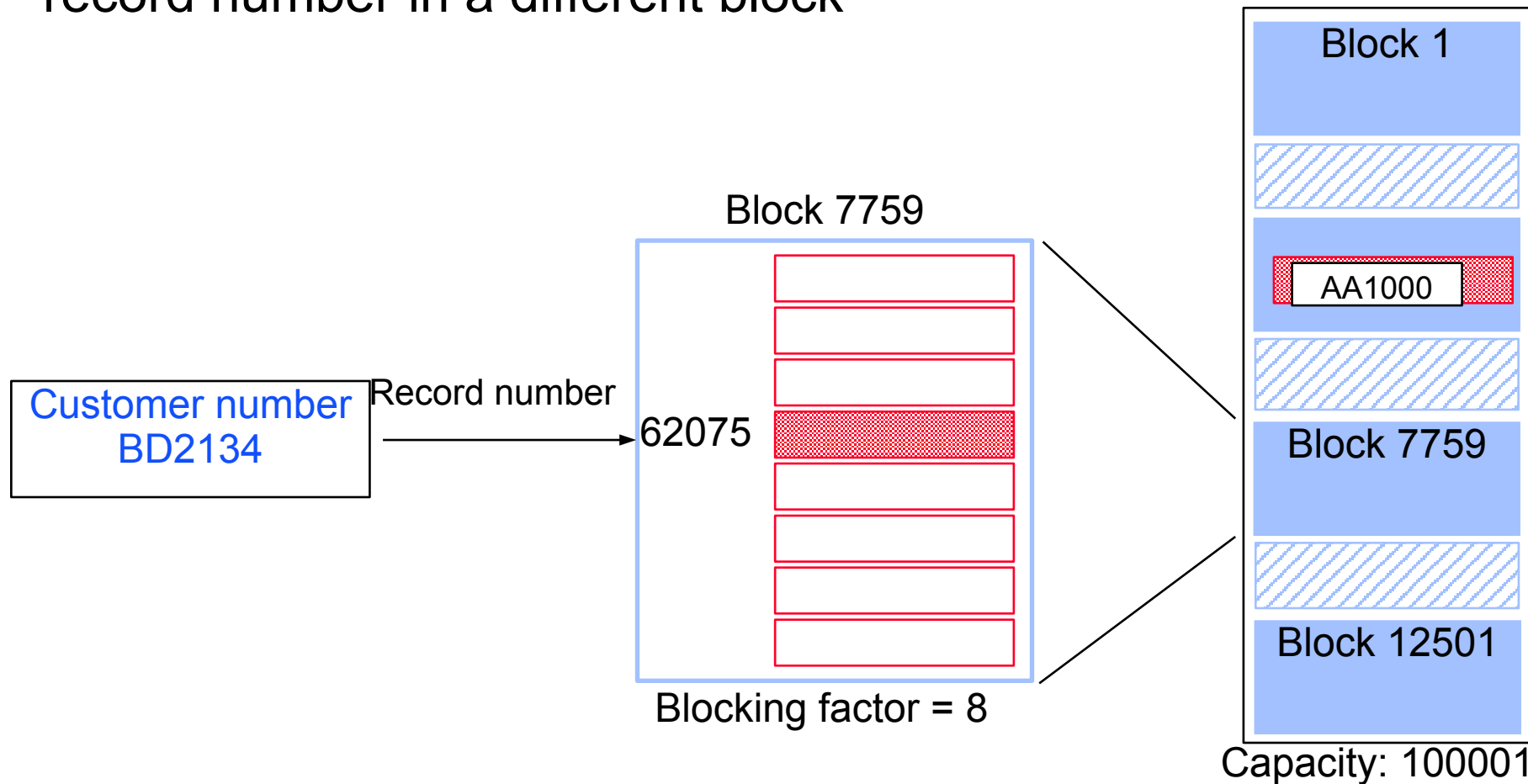
Customer number AA1000 is transformed into a record number



Hashing algorithm (no collision)



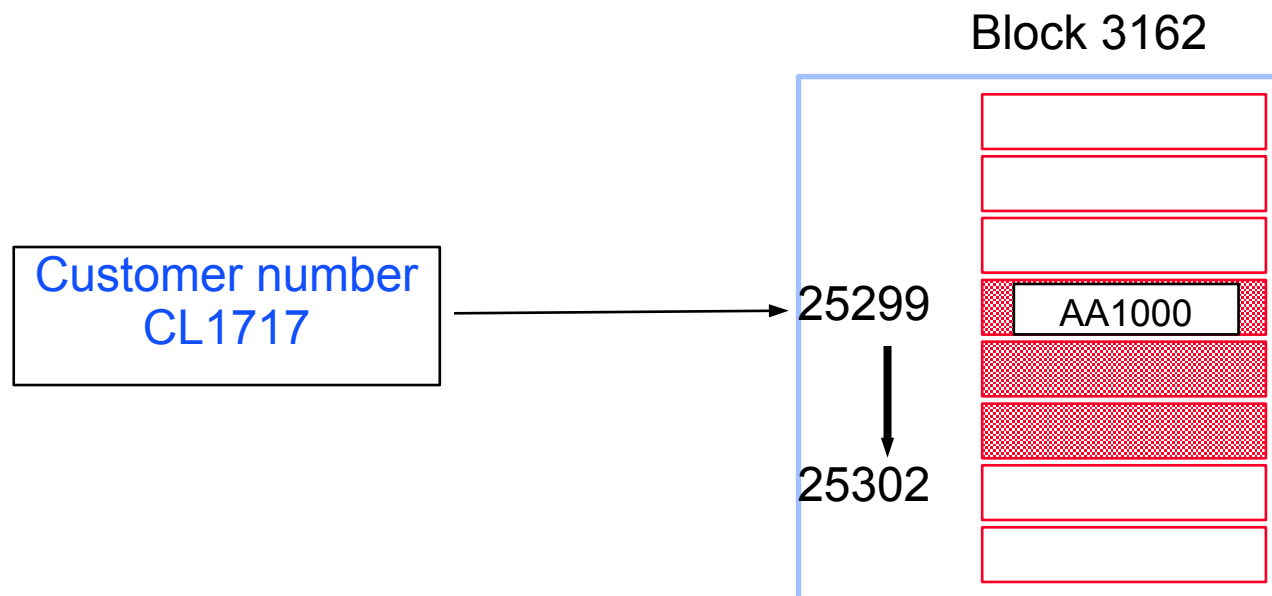
Customer number BD2134 gives a different record number in a different block



Hashing algorithm (collision - same block)



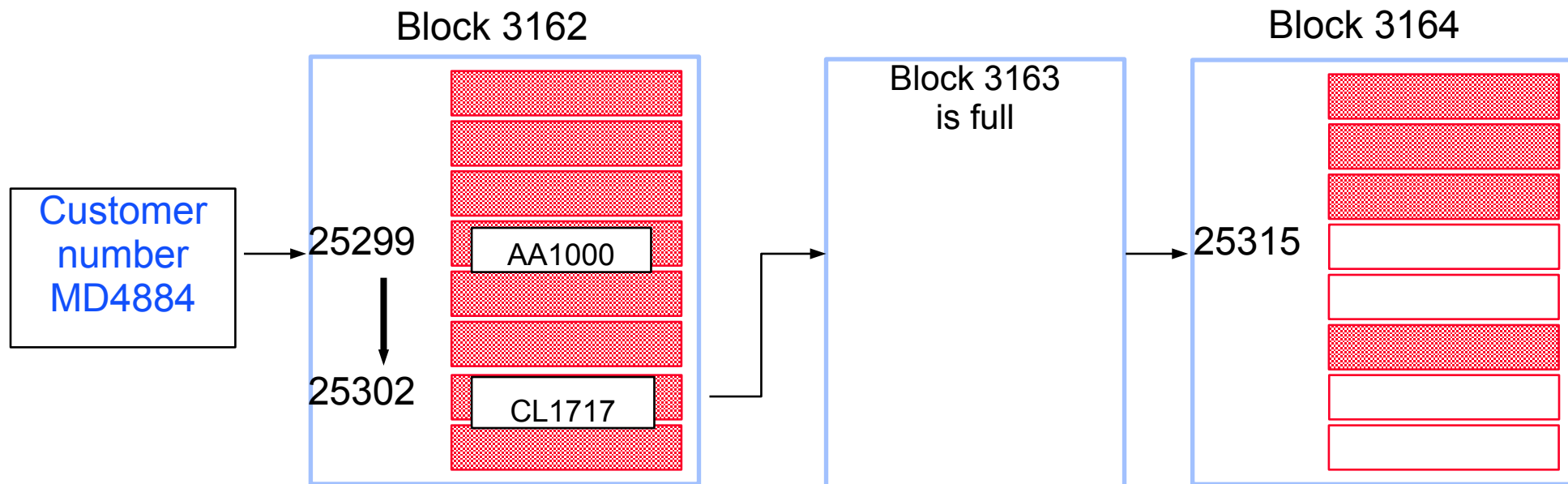
- Customer number CL1717 hashes to the same record number as AA1000 location
- TurboIMAGE tries to find an empty location in the same block. If it finds one, no additional I/O is required.
- CL1717 becomes a secondary entry. Primary and secondary entries are linked using pointers that form a chain.



Hashing algorithm (collision - different block)



- Customer number MD4884 collides with AA1000
- No more room in this block. TurboIMAGE reads the following blocks until it finds a free record location.
- In this case, MD4884 is placed two blocks away, which requires two additional I/Os.

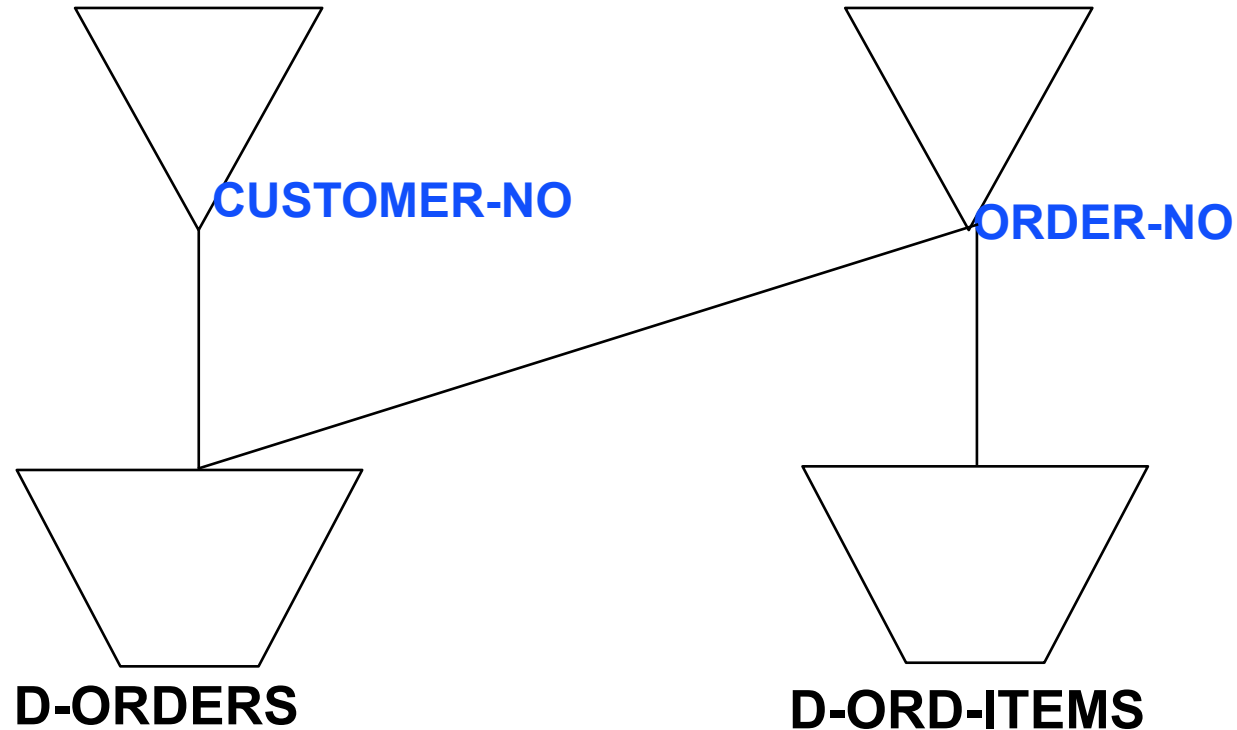


An example TurboIMAGE database



M-CUSTOMER

A-ORDER-NO



HowMessy sample report



HowMessy/XL (Version 2.2.1)
TurboIMAGE/3000 databases

Data Base: STORE.DATA.INVENT
By Robelle Solutions Technology Inc.

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Page: 1

Data Set	Type	Capacity	Entries	Second- Load Factor (Highwater)	Max daries Blks	Blk Fact
M-Customer	Man	248113	178018	71.7%	30.5% 1496	11
A-Order-No	Ato	1266783	768556	60.7%	25.7% 1	70
D-Orders	Det	1000000	768558	76.9%	(851445)	32
D-Ord-Items	Det	4000000	3458511	86.5%	(3470097)	23

Search Field	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs	Elong- ation
Customer-No	32	1.92	0.32	1.00	1.90	90.5%	1.90
Order-No	10	1.35	0.62	1.00	1.00	0.0%	1.00
!Order-No	1	1.00	0	1.00	1.00	0.0%	1.00
S Customer-No	80	14.34	17.76	1.75	9.20	57.2%	5.25
S !Order-No	1604	8.06	35.75	1.36	11.32	72.5%	8.34

HowMessy sample report (master dataset)



HowMessy/XL (Version 2.2.1)
TurboIMAGE/3000 databases

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Data Set	Type	Capacity	Entries	Second- Load Factor	Max daries (Highwater)	Blks	Blk Fact
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D-Orders	Det	1000000	768558	76.9%	(851445)		32
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Search Field	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs	Elong- ation
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Order-No	10	1.35	0.62	1.00	1.00	0.0%	1.00
!Order-No	1	1.00	0	1.00	1.00	0.0%	1.00
S Customer-No	80	14.34	17.76	1.75	9.20	57.2%	5.25
S !Order-No	1604	8.06	35.75	1.36	11.32	72.5%	8.34

Interpreting master datasets lines



☐ Pay attention to the following statistics:

- ☐ High percentage of Secondaries (inefficient hashing)
- ☐ High Maximum Blocks (clustering)
- ☐ High Maximum and Average Chains (inefficient hashing)
- ☐ High Inefficient Pointers (when secondaries exist)
- ☐ High Elongation (when secondaries exist)

Report on m-customer



- The number of Secondaries is not unusually high
- However, there may be problems
 - Records are clustering (high Max Blks)
 - Long synonym chain
 - High percentage of Inefficient Pointers

Data Set	Type	Capacity	Entries	Load Factor	Secon- Max daries Blks (Highwater)	Blk Fact			
M-CUSTOMER	Man	248113	178018	71.7%	<u>30.5%</u> 1496	11			
	Search Field	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs	Elong- ation	
	CUSTOMER-NO	<u>22</u>	1.92	0.32	1.00	1.90	<u>90.5%</u>	1.90	

Report on a-order-no



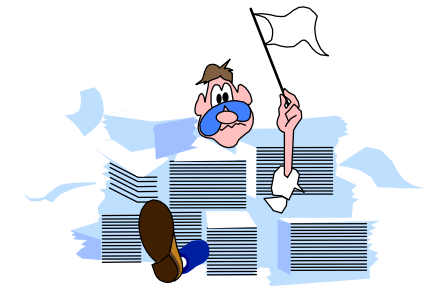
Very tidy dataset

Number of Secondaries is acceptable

Max Blks, Ineff Ptrs and Elongation are at the minimum values, even if the Maximum Chain length is a bit high

Data Set	Type	Capacity	Entries	Load Factor	Secondaries (Highwater)	Max Blks	Blk Fact	Elongation	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs
A-ORDER-NO	Ato	1266783	768556	60.7%	<u>25.7%</u>	<u>1</u>	70	ORDER-NO	<u>10</u>	1.35	0.62	1.00	1.00	<u>0.0%</u> 15
								1.00						

Master dataset solutions



- ❑ Increase capacity to a higher odd number
- ❑ Increase the Blocking Factor
 - ❑ Increase block size
 - ❑ Reduce record size
- ❑ Change binary keys to type X, U, Z, or P
- ❑ Check your database early in the design
- ❑ Use HowMessy on test databases

Implementing a solution: Changing a key data-type



HowMessy/XL (Version 2.2.1)
TurboIMAGE/3000 databases

Data Base: SOMEDB.DATA.SOMEACCT
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Page: 1

Data Set		Type	Capacity	Entries	Second- Max Load daries Blks Factor (Highwater)	Blk Fact
Part-Master	Ato	10000	4305	43.0%	14.2% 16 78	
Part-Loc-Master	Ato	606010	303005	50.0%	<u>86.4%</u> <u>612</u> 67	
Parts-Detail	Det	303030	303005	100.0%	(303005) 63	

Search Field	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs	Elong- ation
Part-No	3	1.17	0.40	1.00	1.67	63.3%	1.67
Part-Loc-No	<u>4305</u>	<u>7.38</u>	73.09	1.07	<u>45.84</u>	66.4%	42.90
Part-No	41065	70.38	666.46	1.86	3.20	3.1%	1.72
!Part-Loc-No	1	1.00	0.00	1.00	1.00	0.0%	1.00

Change key type to byte, same data values



HowMessy/XL (Version 2.2.1)
TurboIMAGE/3000 databases

Data Base: SOMEDB.DATA.SOMEACCT
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Page: 1

Data Set		Type	Capacity	Entries	Second- Max Load daries Factor (Highwater)	Blks	Blk Fact
Part-Master	Ato	10000	4305	43.0%	14.2%	16	78
Part-Loc-Master	Ato	606010	303005	50.0%	<u>27.7%</u>	<u>95</u>	67
Parts-Detail	Det	303030	303005	100.0%	(303005)		63

Search Field	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs	Elong- ation
Part-No	3	1.17	0.40	1.00	1.67	63.3%	1.67
Part-Loc-No	<u>13</u>	<u>1.38</u>	0.77	1.00	<u>1.42</u>	29.3%	1.42
Part-No	41065	70.38	666.46	1.86	3.20	3.1%	1.72
!Part-Loc-No	1	1.00	0.00	1.00	1.00	0.0%	1.00

Change key type to byte, convert data values



HowMessy/XL (Version 2.2.1)
TurboIMAGE/3000 databases

Data Base: SOMEDB.DATA.SOMEACCT
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Page: 1

Data Set		Type	Capacity	Entries	Second- Max Load daries Factor (Highwater)	Blks	Blk Fact
Part-Master	Ato	10000	4305	43.0%	14.2%	16	78
Part-Loc-Master	Ato	606010	303005	50.0%	<u>21.4%</u>	<u>0</u>	56
Parts-Detail	Det	303054	303005	100.0%	(303005)		53

Search Field	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs	Elong- ation
Part-No	3	1.17	0.40	1.00	1.67	63.3%	1.67
Part-Loc-No	<u>6</u>	<u>1.27</u>	0.54	1.00	<u>1.00</u>	0.0%	1.00
Part-No	41065	70.38	666.46	2.07	3.43	3.4%	1.66
!Part-Loc-No	1	1.00	0.00	1.00	1.00	0.0%	1.00

HowMessy sample report (detail dataset)



HowMessy/XL (Version 2.2.1)
for TurboIMAGE/3000 databases

Data Base: STORE.DATA.INVENT
By Robelle Solutions Technology Inc.

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Page: 1

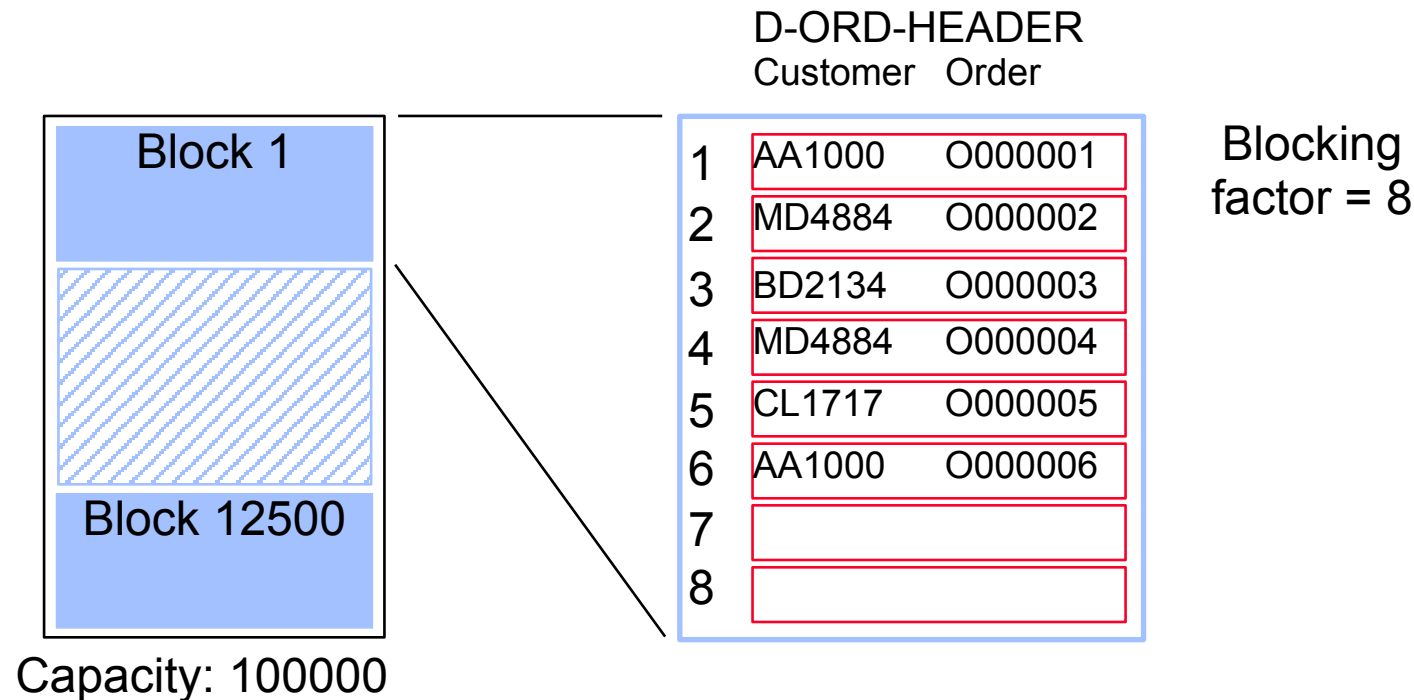
Data Set	Type	Capacity	Entries	Load Factor	Secon- Max daries Blks (Highwater)	Blk Fact
M-CUSTOMER	Man	248113	178018	71.7%	30.5% 1496	1
A-ORDER-NO	Ato	126673	768556	60.7%	25.7% 1	70
D-ORDERS	Det	1000000	768556	76.9%	(851445)	12
D-ORD-ITEMS	Det	4000000	3458511	86.5%	(3470097)	23

Search Field	Max Chain	Ave Chain	Std Dev	Expd Blocks	Avg Blocks	Ineff Ptrs	Elong- ation
Customer-No	22	1.92	0.32	1.00	1.90	90.5%	1.90
Order-No	10	1.35	0.62	1.00	1.00	0.0%	1.00
!Order-No	1	1.00	0	1.00	1.00	0.0%	1.00
S Customer-No	80	14.34	17.76	1.75	9.20	57.2%	5.25
S !Order-No	1604	8.06	35.75	1.36	11.32	72.5%	8.34

Empty detail dataset



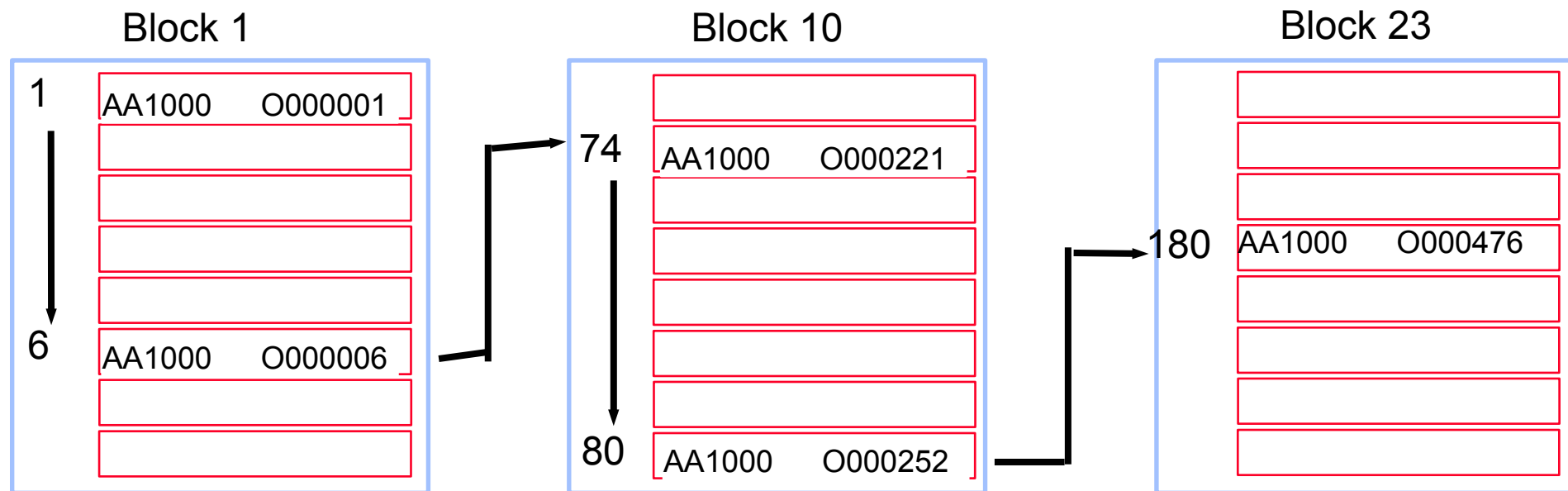
- Records are stored in the order they are created, starting from record 1
- Records for the same customer are linked together using pointers to form a chain
- Chains are linked to the corresponding master entry



Detail chains get scattered



Over time, records for the same customer are scattered over multiple blocks

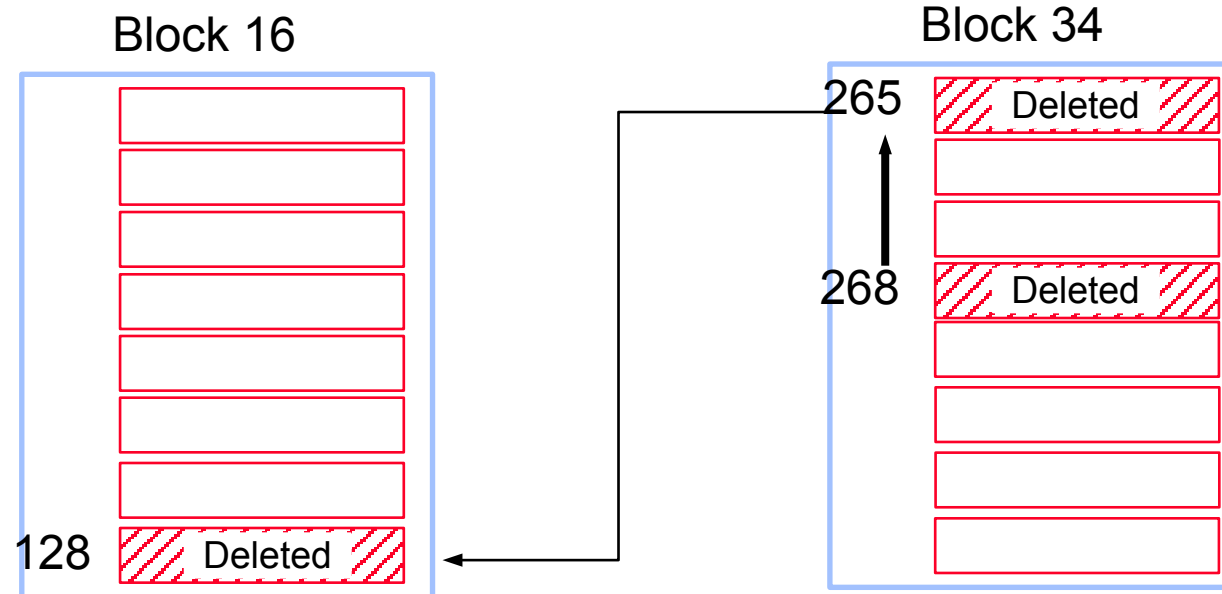


Delete chain



Deleted records are linked together

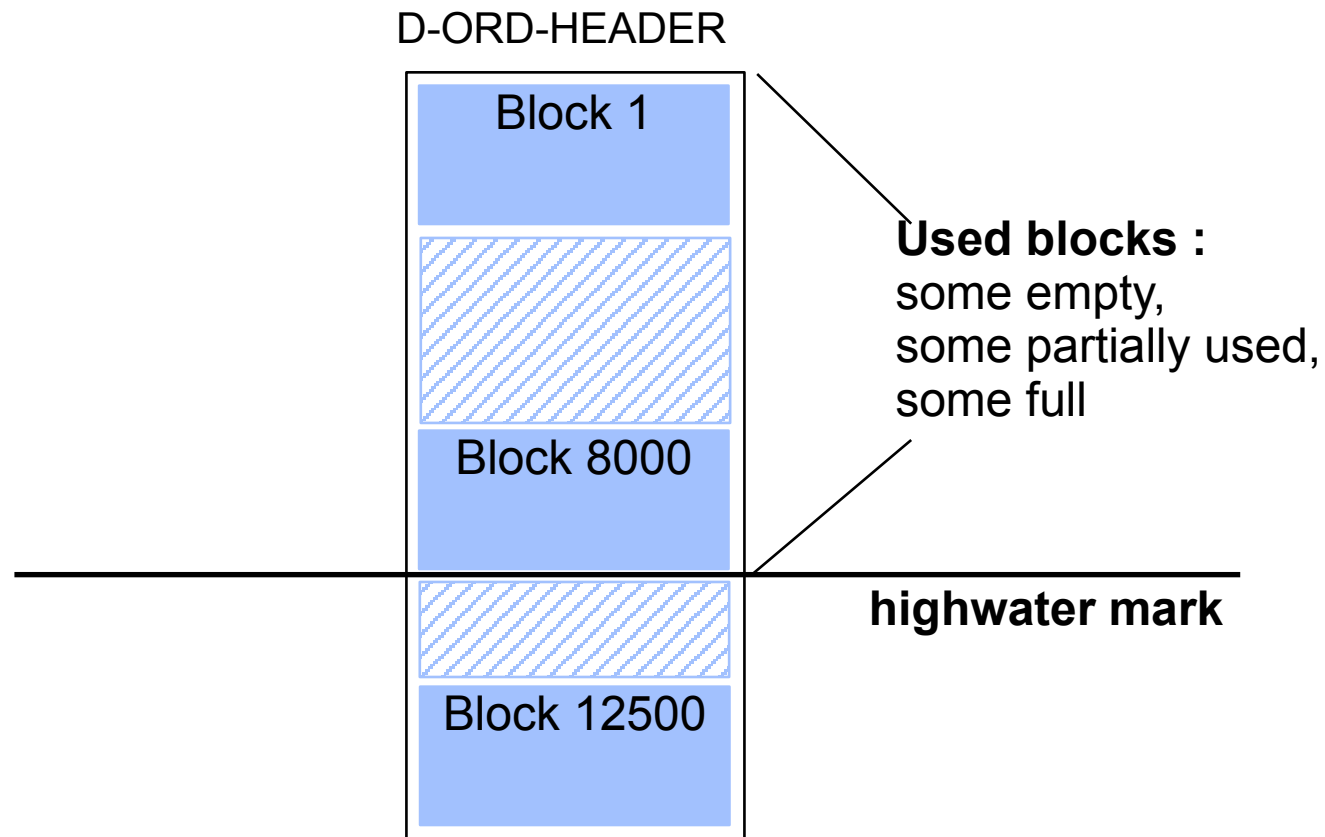
TurboIMAGE reuses the locations in the Delete chain, if there are any



Highwater mark



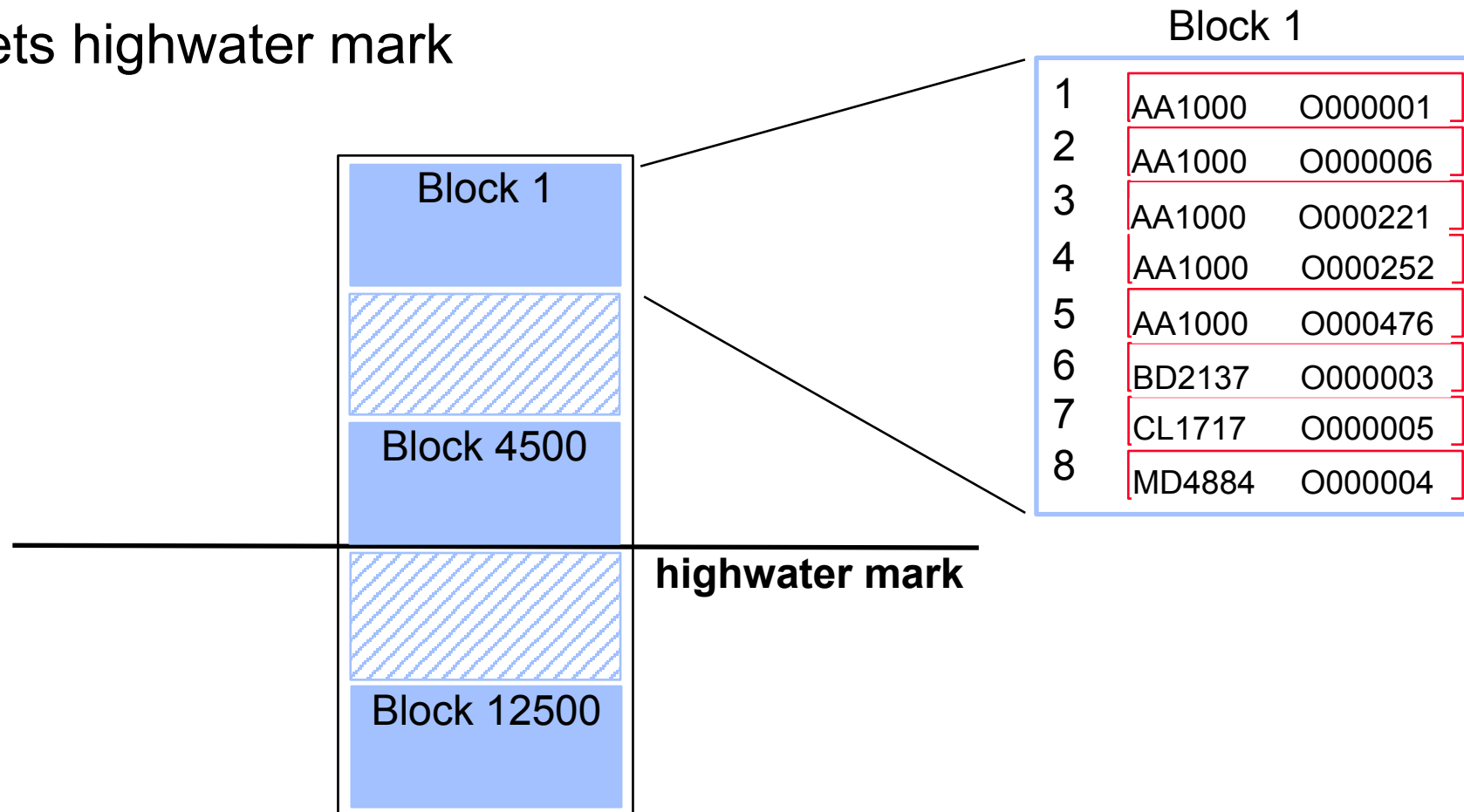
- Indicates highest record location used so far
- Serial reads scan the dataset up to the highwater mark



Repacking a detail dataset



- Groups records along primary path
- Removes Delete chain (no holes)
- Resets highwater mark



Interpreting detail dataset lines



☐ Pay attention to the following statistics:

- ☐ Load Factor approaching 100% (dataset full)
- ☐ Primary path (large Average Chain and often accessed)
- ☐ High Average Chain and low Standard Deviation, especially with a sorted path (Is path really needed?)
- ☐ High Inefficient Pointers (entries in chain not consecutive)
- ☐ High Elongation (entries in chain not consecutive)

Report on d-orders



- Primary path should be on customer-no, not on order-no
- Highwater mark is high
- Repack along new primary path regularly

Data Set	Type	Capacity	Entries	Load Factor	Secon-Max daries Blks (Highwater)	Blk Fact			
D-ORDERS	Det	1000000	768556	76.9%	(<u>851445</u>)	12			
		Max	Ave	Std	Expd	Avg	Ineff	Elong-	
	Search Field	Chain	Chain	Dev	Blocks	Blocks	Ptrs	ation	
	!ORDER-NO	1	1.00	0	1.00	1.00	0.0%	1.00	
S	CUSTOMER-NO		<u>80</u>	<u>14.34</u>	17.76	1.75	9.20	<u>57.2%</u>	
5.25									

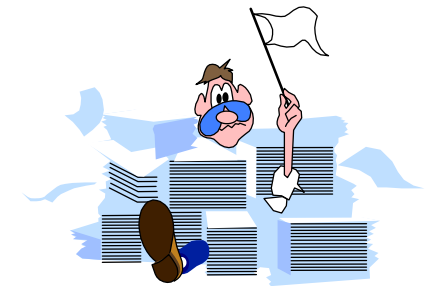
Report on d-ord-items



- NUL Inefficient Pointers and Elongation are high
- NUL Highwater mark is fairly high
- NUL Repack the dataset regularly
- NUL Is the sorted path really needed?

Data Set	Type	Capacity	Entries	Load Factor	Secon- Max daries Blks (Highwater)	Blk Fact		
D-ORD-ITEMS	Det	4000000	3458511	86.5%	(<u>3470097</u>)	23		
		Max	Ave	Std	Expd	Avg	Ineff	Elong-
	Search Field	Chain	Chain	Dev	Blocks	Blocks	Ptrs	ation
	S !ORDER-NO	1604	8.06	35.75	1.36	<u>11.32</u>	<u>72.5</u>	8.34

Detail dataset solutions



- ❑ Assign the primary path correctly; select a search item with an Average Chain length > 1 that is accessed most often
- ❑ Repack datasets along the primary path regularly
- ❑ Increase the Blocking Factor
 - ❑ Increase block size
 - ❑ Reduce record size
- ❑ Understand sorted paths
- ❑ Check your databases early in the design; use HowMessy on test databases

Minimum number of disc I/Os



Intrinsic

Disc I/Os

DBGET	1
DBFIND	1
DBBEGIN	1
DBEND	1
DBUPDATE	1 (non-critical item)
DBUPDATE	13 (critical item)
DBPUT	3 [+ (4 x #paths, if detail)]
DBDELETE	2 [+ (4 x #paths, if detail)]

Serial reads:

Master
Detail

Capacity / Blocking factor
entries / Blocking factor

Estimating response time



NUL Deleting 100,000 records from a detail dataset with two paths would take:

NUL $2 + (4 \times 2 \text{ paths}) = 10 \text{ I/Os per record}$

NUL $100,000 \text{ records} \times 10 \text{ I/Os per record} = 1,000,000 \text{ I/Os}$

NUL Classic: around 25 I/Os per second

NUL $1,000,000 \text{ I/Os} / 25 = 40,000 \text{ seconds}$

NUL $40,000 \text{ seconds} / 3600 = 11.1 \text{ hours}$

NUL iX: around 40 I/Os per second

NUL $1,000,000 \text{ I/Os} / 40 = 25,000 \text{ seconds}$

NUL $25,000 \text{ seconds} / 3600 = 6.9 \text{ hours}$

Automating HowMessy analysis



- Recent version of HowMessy creates a self-describing file with these statistics
- Process the file with generic tools (Suprtool, AskPlus) or custom programs (COBOL, 4GL), and produce custom reports
- Send messages to database administrators
- Write “smart” job to fix databases without user intervention

Processing Loadfile with Suprtool



■ Datasets more than 80% full

```
>input loadfile
>if loadfactor > 80
>ext database, dataset, datasettype, loadfactor
>list standard
```

■ Only one address per customer

```
>input loadfile
>if dataset = "D-ADDRESSES" and &
maxchain > 1
```

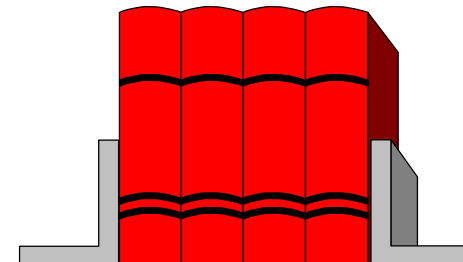
References

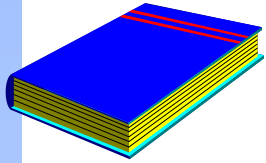


 The TurboIMAGE/3000 Handbook (Chapter 23)

 Available for \$ 49.95 from:

WORDWARE
P.O. Box 14300
Seattle, WA 98114





Summary



- ❑ TurboIMAGE databases become messy over time, especially if they are active
- ❑ HowMessy and DBLOADNG let you analyze the database's efficiency
- ❑ You should have some knowledge of the internal workings of TurboIMAGE
- ❑ Monitor your databases regularly