

**HowToCode7**

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# Chapter 1

## HowToCode7

### 1.1 HowToCode: Optimising

Optimising your Code

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Everyone wants their code to run as fast as possible, so here are some speed-up tricks for you:

- 1 68000 Optimisation
- 2 68020 Optimisation
- 3 Blitter Speed Optimisation
- 4 General Speed-Up Notes

### 1.2 68000 Optimisation

68000 optimisation

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Some notes added by CJ

NOTE! Not all these optimisations can be automatically applied. Make sure they will not affect other areas in your code!

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Original	Possible optimisation	Examples/notes
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STANDARD WELL-KNOWN optimisations

RULE: use Quick-type/Short branch! Use INLINE subroutines if they are small!

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BRA/BSR xx	BRA.s/BSR.s xx	if xx is close to PC
MOVE.X #0	CLR.X/MOVEQ/SUBA.X	move.l #0,count -> clr.l count
move.l #0,d0	-> moveq #0,d0	
move.l #0,a0	-> sub.l a0,a0	

---

CLR.L Dx	MOVEQ #0,Dx	
CMP #0	TST	
MOVE.L #nn,dx	MOVEQ #nn,dx	possible if -128<=nn<=127
ADD.X #nn	ADDQ.X #nn	possible if 1<=nn<=8
SUB.X #nn	SUBQ.X #nn	same...
JMP/JSR xx	BRA/BSR xx	possible if xx is close to PC
JSR xx;RTS	JMP xx	save a RTS
BSR xx;RTS	BRA xx	same... (assuming routine doesn't rely on anything in the stack)

LSL/ASL #1/2,xx ADD xx,xx [ADD xx,xx] lsl #2,d0 -> 2 times add d0,d0

MULU #yy,xx where yy is a power of 2, 2..256  
 LSL/ASL #1-8,xx mulu #2,d0 -> asl #1,d0 -> add d0,d0  
 BEWARE: STATUS FLAGS ARE "WRONG"

DIVU #yy,xx where yy is a power of 2, 2..256  
 LSR/ASR #.. SWAP divu #16,d0 -> lsr #4,d0  
 BEWARE: STATUS FLAGS ARE "WRONG",  
 AND HIGHWORD IS NOT THE REMAINDER.

## ADDRESS-RELATED OPTIMISATIONS

RULE: use short addressing/quick adds!

-----

MOVEA.L #nn	MOVEA.W #nn	Movea is "sign-extending" thus possible if 0<=nn<=\$7fff
ADDA.X #nn	LEA nn()	adda.l #800,a0 -> lea 800(a0),a0 possible if -\$8000<=nn<=\$7fff
LEA nn()	ADDQ.W #nn	lea 6(a0),a0 -> addq.w #6,a0 possible if 1<=nn<=8
\$0000nnnn.l	\$nnnn.w	move.l 4,a6 -> move.l 4.w,a6 possible if 0<=nnnn<=\$7fff (nnnn is SIGN EXTENDED to LONG!)
MOVE.L #xx,Ay	LEA xx,Ay	try xx(PC) with the LEA

MOVE.L Ax,Ay;  
 ADD #nnnn,Ay LEA nnnn(Ax),Ay copy&add in one

## OFFSET-RELATED OPTIMISATIONS

RULE: use PC-relative addressing or basereg addressing!  
 put your code&data in ONE segment if possible!

-----

MOVE.X nnnn	MOVE.X nnnn(pc)	lea copper,a0 -> lea copper(pc),a0..
LEA nnnn	LEA nnnn(pc)	...possible if nnnn is close to PC
(Ax,Dx.l)	(Ax,Dx.w)	possible if 0<=Dx<=\$7fff

If PC-relative doesn't work, use Ax as a pointer to your data block.  
Use indirect addressing to get to your data: `move.l Data1-Base(Ax),Dx` etc.

#### TRICKY OPTIMISATIONS

```
-----
BSET #xx,yy      ORI.W #2^xx,yy      0<=xx<=15
BCLR #xx,yy      ANDI.W #~(2^xx),yy  "
BCHG #xx,yy      EORI.W #2^xx,yy     "
BTST #xx,yy      ANDI.W #2^xx,yy     "
                Best improvement if yy=a data reg.
                BEWARE: STATUS FLAGS ARE "WRONG".
```

#### SILLY OPTIMISATIONS (FOR OPTIMISING COMPILER OUTPUTS ETC)

```
-----
MOVEM (one reg.)  MOVE.L              movem d0,-(sp) -> move.l d0,-(sp)

MOVE xx,-(sp)     PEA xx              possible if xx=(Ax) or constant.

0 (Ax)            (Ax)

MULU/MULS #0      CLR.L              moveq #0,Dx with data-registers.

MULU #1,xx        SWAP CLR SWAP      high word is cleared with mulu #1
MULS #1,xx        SWAP CLR SWAP EXT.L see MULU, and sign extended.
                BEWARE: STATUS FLAGS ARE "WRONG"
```

#### LOOP OPTIMISATION.

Example: imagine you want to eor 4096 bytes beginning at (a0).  
Solution one:

```
        move.w    #4096-1,d7
.1      eori.b    d0,(a0)+
        dbra      d7,.1
```

Consider the loop from above. 4096 times a `eor.b` and a `dbra` takes time.  
What do you think about this:

```
        move.w    #4096/4-1,d7
.1      eor.l     d0,(a0)+      ; d0 contains byte repeated 4 times
        dbra      d7,.1
```

Eors 4096 bytes too! But only needs 1024 `eor.l/dbra`s.  
Yeah, I hear you smart guys cry: what about 1024 `eor.l` without any loop?!  
Right, that IS the fastest solution, but is VERY memory consuming (2 Kb).  
Instead, join a loop and a few `eor.l`:

```
        move      #4096/4/4-1,d7
.1      eor.l     d0,(a0)+
        eor.l     d0,(a0)+
        eor.l     d0,(a0)+
        eor.l     d0,(a0)+
        dbra      d7,.1
```

This is faster than the loop before. I think about 8 or 16 `eor.l`'s is just fine, depending on the size of the mem to be handled (and the wanted

speed!). Also, mind the cache on 68020+ processors, the loop code must be small enough to fit in it for highest speeds. Try to do as much as possible within one loop (but considering the text above) instead of a few loops after each other.

#### MEMORY CLEARING/FILLING.

-----  
A common problem is how to clear or fill some memory in a short time. If it is CHIP-MEMORY, use the blitter (only D-channel, see below). In this case you can still do other things with your 680x0 while the blitter is busy erasing. If it is FAST-MEMORY, you can use the method from above, with `clr.l` instead of `eor.l`, but there is a much faster way:

```

    move.l    sp,TempSp
    lea       MemEnd,sp
    moveq     #0,d0
;...for all 7 data regs...
    moveq     #0,d7
    move.l    d0,a0
;...for 6 address regs...
    move.l    d0,a6

```

After this, ONE instruction can clear 60 bytes of memory (15\*4):

```

    movem.l   d0-d7/a0-a6,-(sp)    ;wham!

```

Now, repeat this instruction as often as required to erase the memory. (memsize/60 times). You may need an additional `movem.l` to erase the last few bytes. Get `sp(=a7)` back at the end with (guess..):

```

    move.l    TempSp,sp

```

If you are low on mem, put a few `movem.l` in a loop. But, now you need a loop-counter register, so you'll only clear 56 bytes in one `movem.l`.

In the case of CHIP memory, you can use both the blitter and the processor simultaneously to clear much CHIP mem in a VERY short time... It takes some experimentation to find the best sizes to clear with the blitter and with the processor.

BUT, ALWAYS USE A `WaitBlit()` AFTER CLEARING SIMULTANEOUSLY, even if you think you know that the blitter is finished before your processor is.

## 1.3 General Speed-Up Notes

- When optimising programs first try to find the time-critical parts (inner loops, interrupt code, often called procedures etc.) In most cases 10% of the code is responsible for 90% of the execution time. Don't waste time doing needless optimising on startup and exit code when it's only called once!
  - Often it is better not to set `BLTPRI` in `DMACON` (#10 in `$dff09a`) as this can keep your processor from calculating things while the blitter is busy.
  - Use as many registers as possible! Store values in registers rather
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than in memory, it's much faster!

- DON'T put your parameters on the stack before calling a routine! Instead, put them in registers!
  - If you have enough memory, try to remove as many MULU/S and DIVU/S as possible by pre-calculating a multiplication or division table, and reading values from it, or rewrite multiply/divide code with simpler instructions if possible (eg ADD, LSR, etc.)
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